

BoundedVersionedString

- A VersionedString with a bounded capacity

- Abstraction Function:

$$A: BVS \rightarrow \{\phi\} \cup S \cup S \times S \cup \dots \cup S^{capacity}$$

where:

- $BVS = \{vs \mid vs \text{ is a BoundedVersionedString state}\}$
- $S = \{s \mid s \text{ is a String}\}$

$$vs \in BVS \xrightarrow{A} \begin{cases} \phi \\ s \in S \\ (s_1, s_2) \in S \times S \\ (s_1, s_2, s_3) \in S \times S \times S \end{cases}$$

Software 1

Recitation No. 5 (Class Verification)

The Contract Pre/Post Conditions

- BoundedVersionedString(int capacity):

- Requires:
capacity > 0
- Ensures:
nothing



Method Specification

```
public class BoundedVersionedString {
    public BoundedVersionedString(int capacity) {...}
    public void add(String s) {...}
    public int length() {...}
    public String getLastVersion() {...}
    public String getVersion(int i) {...}
}
```

The Contract Pre/Post Conditions (cont.)

- int length():

- Requires:
nothing
- Ensures:

$$A(new) = A(old)$$

$$\text{returned value} = \begin{cases} 0 & \text{if } A(old) = \phi \\ k & \text{if } A(old) = (s_1, \dots, s_{k \geq 1}) \end{cases}$$

Since length() is a query

The Contract Pre/Post Conditions (cont.)

- void add(String s):

- Requires:
s != null
- Ensures:

$$A(old) = \phi \Rightarrow A(new) = (s)$$

$$A(old) = (s_1, s_2, \dots, s_{k \geq 1}) \Rightarrow A(new) = \begin{cases} (s_1, s_2, \dots, s_k, s) & \text{if } k < \text{capacity} \\ (s_2, s_3, \dots, s_k, s) & \text{if } k = \text{capacity} \end{cases}$$

The Contract

Pre/Post Conditions (cont.)

■ String getLastVersion():

- Requires:

$$A(\text{this}) \neq \phi$$

- Ensures:

$$A(\text{new}) = A(\text{old})$$

$$A(\text{old}) = (s_1, \dots, s_k) \Rightarrow \text{returned value} = s_k$$

The Contract

Pre/Post Conditions (cont.)

■ String getVersion(int i):

- Requires:

$$A(\text{this}) = (s_1, s_2, \dots, s_{k \geq i \geq 1})$$

- Ensures:

$$A(\text{new}) = A(\text{old})$$

$$A(\text{old}) = (s_1, \dots, s_{k \geq i \geq 1}) \Rightarrow \text{returned value} = s_i$$

Implementation (cont.)

Fields

// Stores the remembered versions
private String[] versions ;

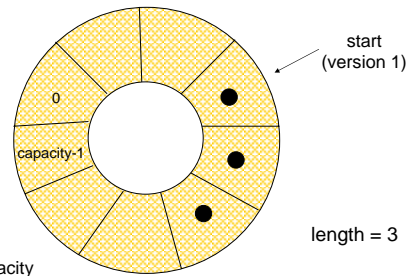
// The maximum number of versions to remember
private int capacity;

// The actual number of remembered versions
private int length = 0;

// The position of the oldest remembered version
private int start = 0;

Implementation

- Based on a circular array



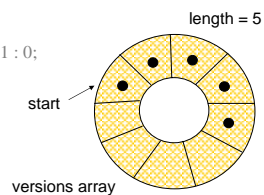
Implementation (cont.)

Command

```
public void add(String s) {
    int pos = (start + length) % capacity;
    versions[pos] = s;

    if ((length > 0) && (start == pos)) {
        start = (start+1) % capacity;
    }

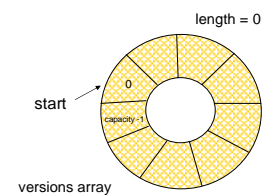
    length += (length < capacity) ? 1 : 0;
}
```



Implementation (cont.)

Constructor

```
public BoundedVersionString(int capacity) {
    this.capacity = capacity;
    versions = new String[capacity];
}
```



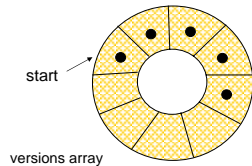
Abstraction Function Definition

$A: BVS \rightarrow \{\emptyset\} \cup S \cup S \times S \cup \dots \cup S^{capacity}$

where:

- $BVS = \{vs \mid vs \text{ is a BoundedVersionedString state}\}$
- $S = \{s \mid s \text{ is a String}\}$

$A(this) = (versions[start], \dots, versions[(start + length - 1) \% capacity])$



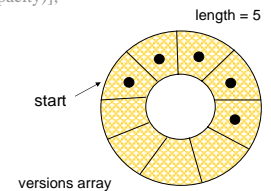
Implementation (cont.)

Queries

```
public int length() {
    return length;
}

public String getVersion(int i) {
    return versions[(start+i-1)%capacity];
}

public String getLastVersion() {
    return getVersion(length());
}
```



Class Verification

■ Invariant Verification in initial state:

- $capacity > 0$ (construction precondition) ■ 1
- $start = 0 < capacity$ (initial state, 1) ■ 2
- $length = 0 < capacity, start = 0$ (initial state) ■ 3
- $A(this) = (versions[0], \dots, versions[-1]) = \emptyset$ ■ 4
- $length = 0$

Representation Invariant

1. $capacity > 0$
2. $0 \leq start < capacity$
3. $length < capacity \Rightarrow start = 0$
4. $length = \begin{cases} 0 & \text{if } A(this) = \emptyset \\ k & \text{if } A(this) = (s_1, \dots, s_{k \geq 1}) \end{cases}$

Class Verification (cont.)

■ Method Verification:

- Assume: The invariant and precondition are satisfied **before** the call
- Prove: The invariant and postcondition are satisfied **after** the call

■ Other Code Verification:

- Prove: No other code in the program change the invariant.