

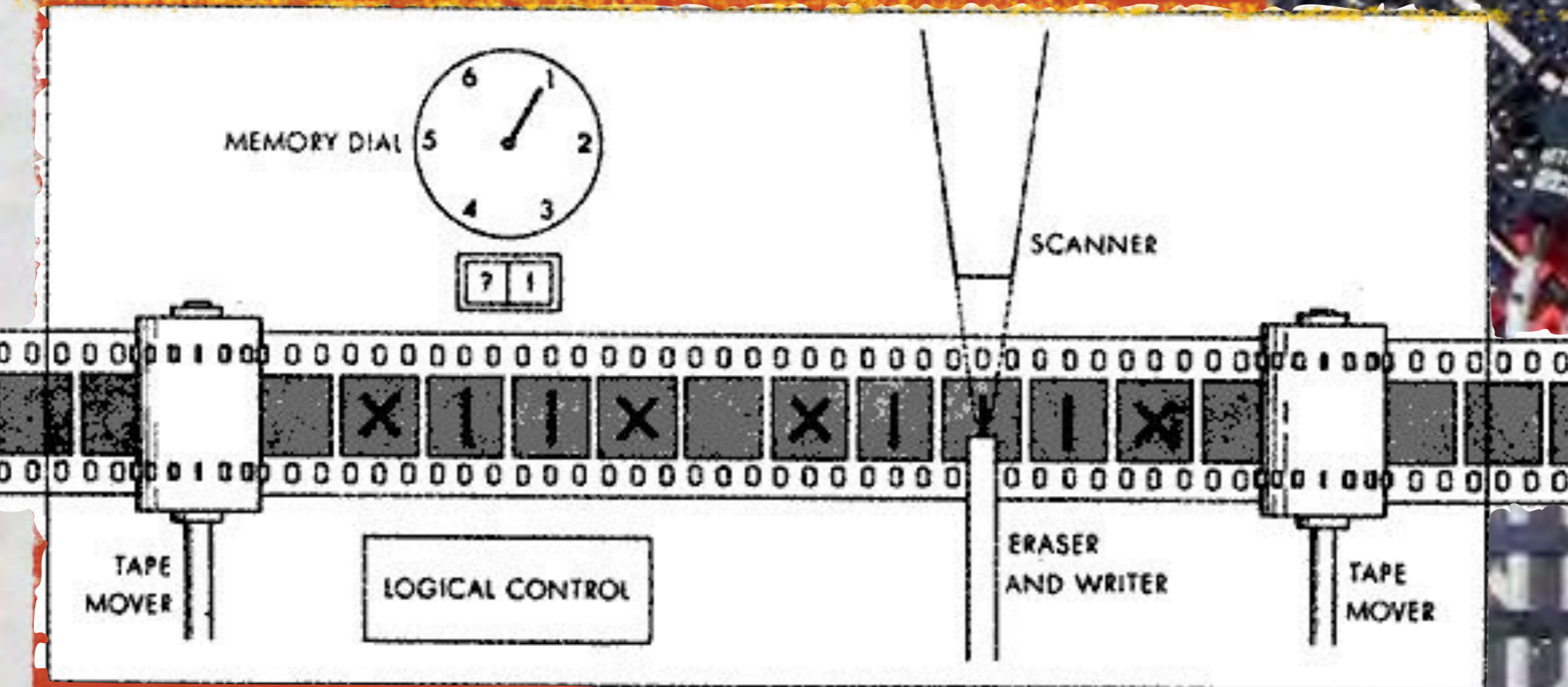
Analog Models

Alan Turing (1936)

A man provided with paper, pencil, and rubber, and subject to strict discipline, is in effect a universal machine



Turing's Machines



Turing [1948]

- The property of being 'discrete' is only an advantage for the theoretical investigator, and serves no evolutionary purpose, so we could not expect Nature to assist us by producing truly 'discrete' brains.

Turing [1948]

- All machinery can be regarded as continuous, but when it is possible to regard it as discrete it is usually best to do so.

Turing's Premises

- Sequential (discrete) symbol manipulation
- Deterministic
- Finite internal states
- Finite symbol space
- Finite observability and local action
- Linear external memory suffices

Digital vs. Analog



Turing [1950]

- It is true that a discrete-state machine must be different from a continuous machine. But if we adhere to the conditions of the imitation game, the interrogator will not be able to take any advantage of this difference.

Analog

- Analog space
- Analog time



Analog Space

- Uncountably many possible values



Euclid (c. -300)

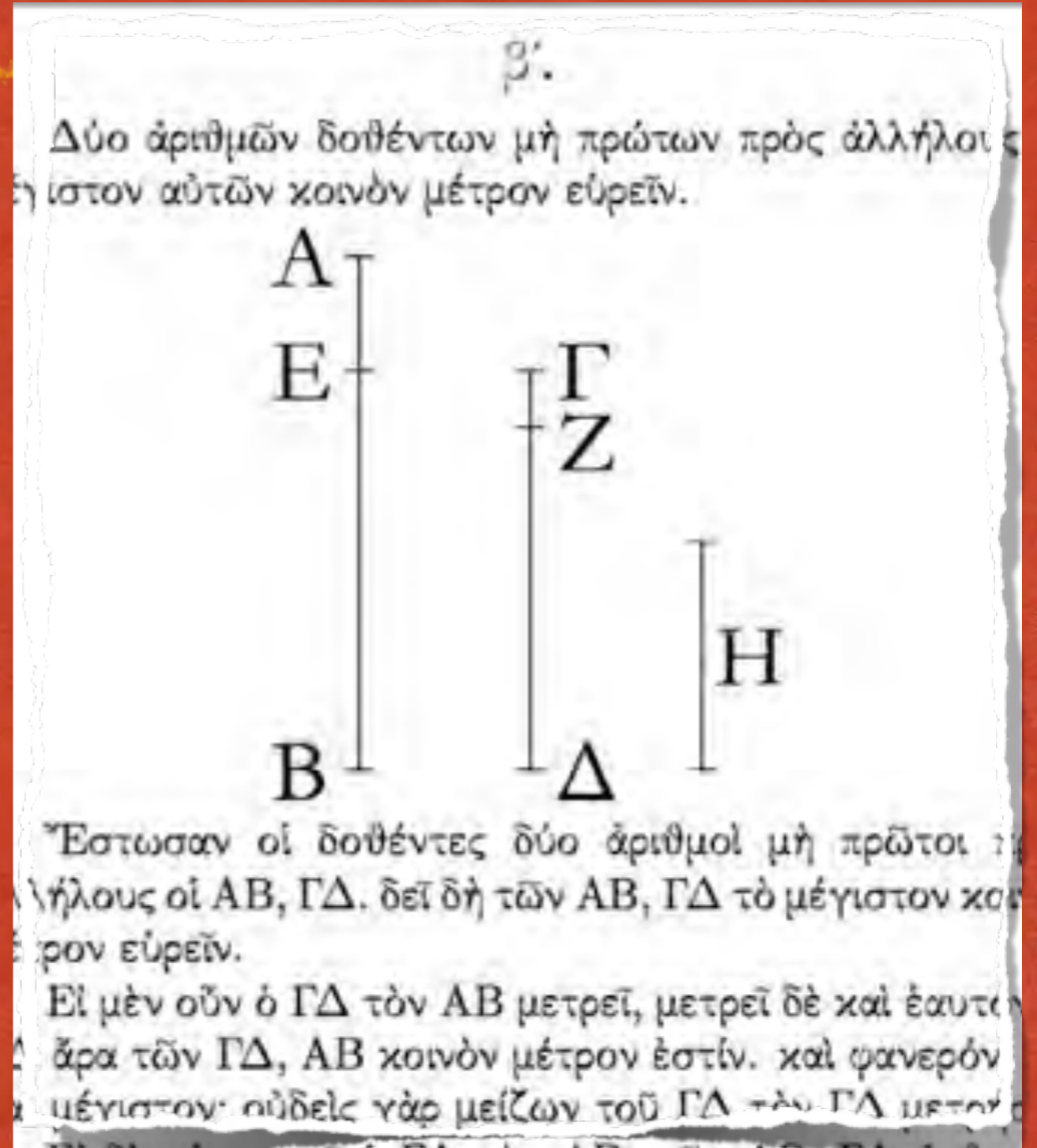


Euclid's GCD algorithm was formulated geometrically: Find common measure for 2 lines.

Used repeated subtraction of the shorter segment from the longer.

Euclid's Elements

- Finitely describable
— in terms of basic compass operations



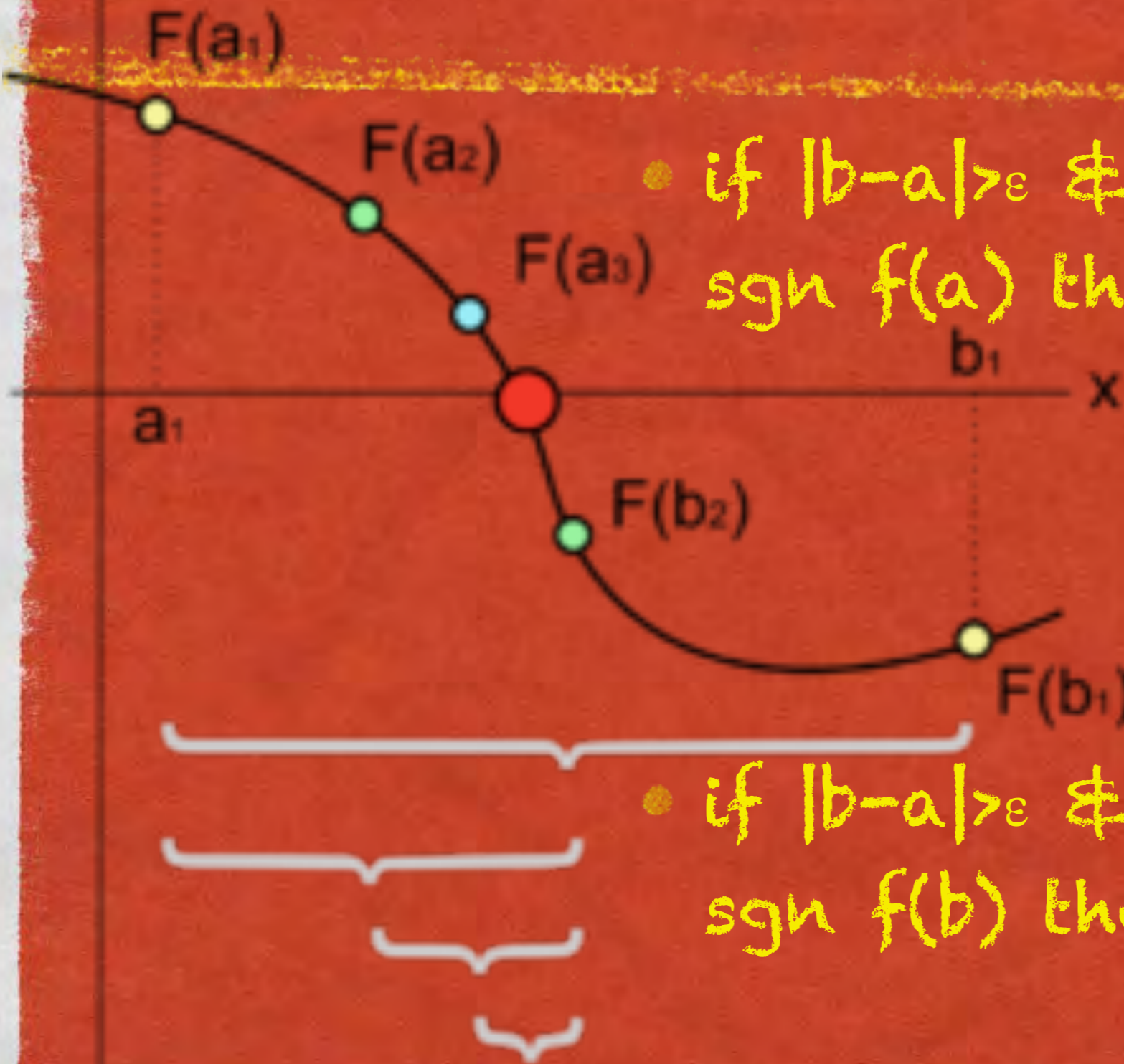
Turing [1950]

- A small error in the information about the size of a nervous impulse impinging on a neuron, may make a large difference to the size of the outgoing impulse.



$F(x)$

Bisection Search



- if $|b-a| > \epsilon$ & $\text{sgn } f((a+b)/2) = \text{sgn } f(a)$ then $a := (a+b)/2$

- if $|b-a| > \epsilon$ & $\text{sgn } f((a+b)/2) = \text{sgn } f(b)$ then $b := (a+b)/2$

Bisection Search (Saul Gorn)

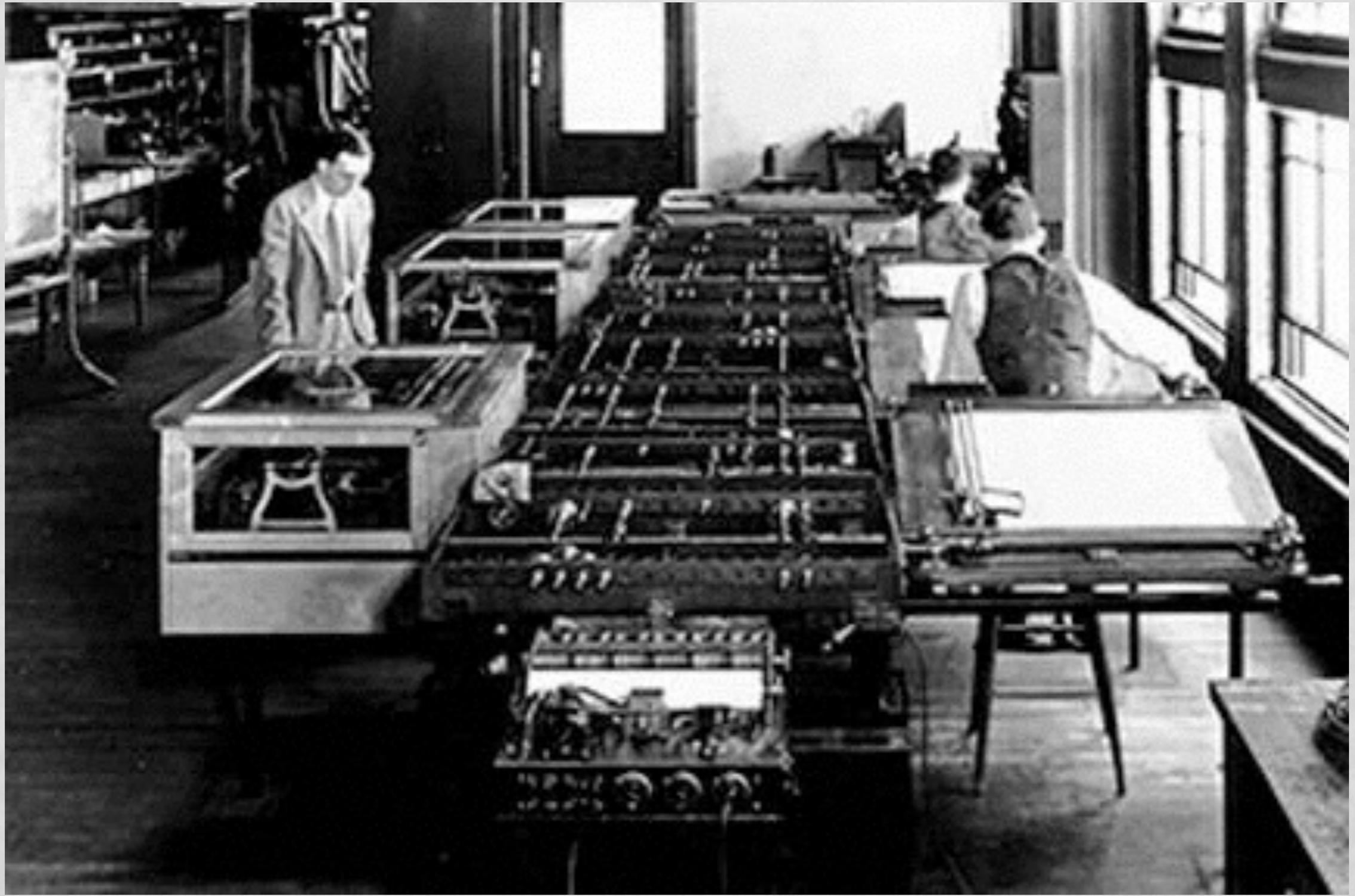
- Although [this procedure is] among the slowest, it is applicable to any continuous function. The fact that no differentiability conditions have to be checked makes it ... an 'old work-horse'.



Continuous Space

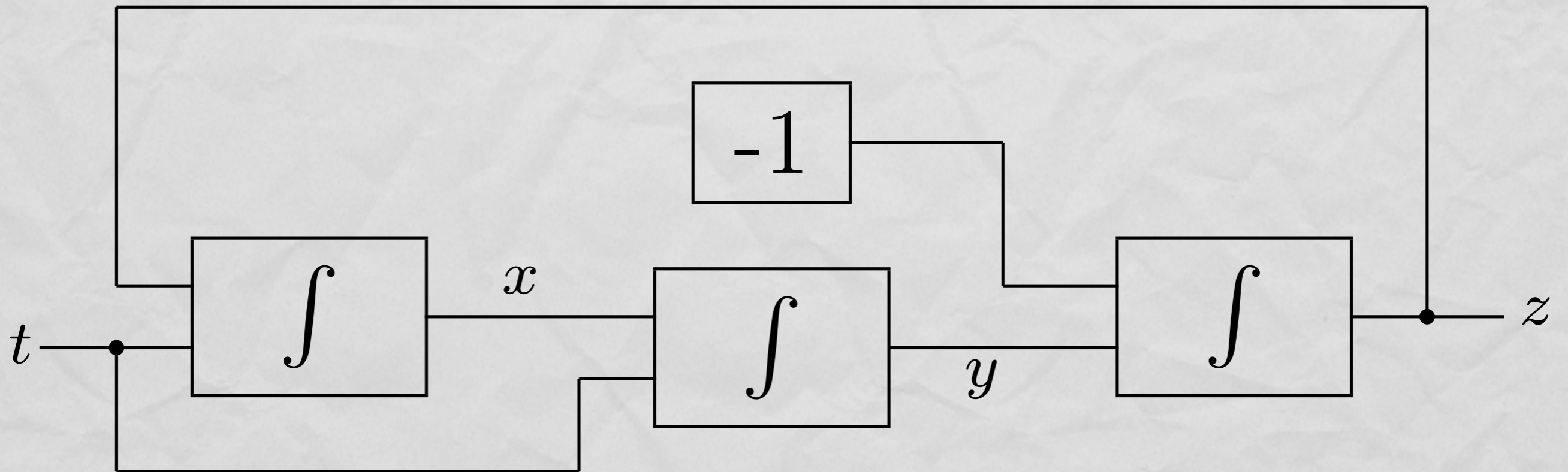
- Idealized
- Computable reals
- Intervals
- Arbitrary precision

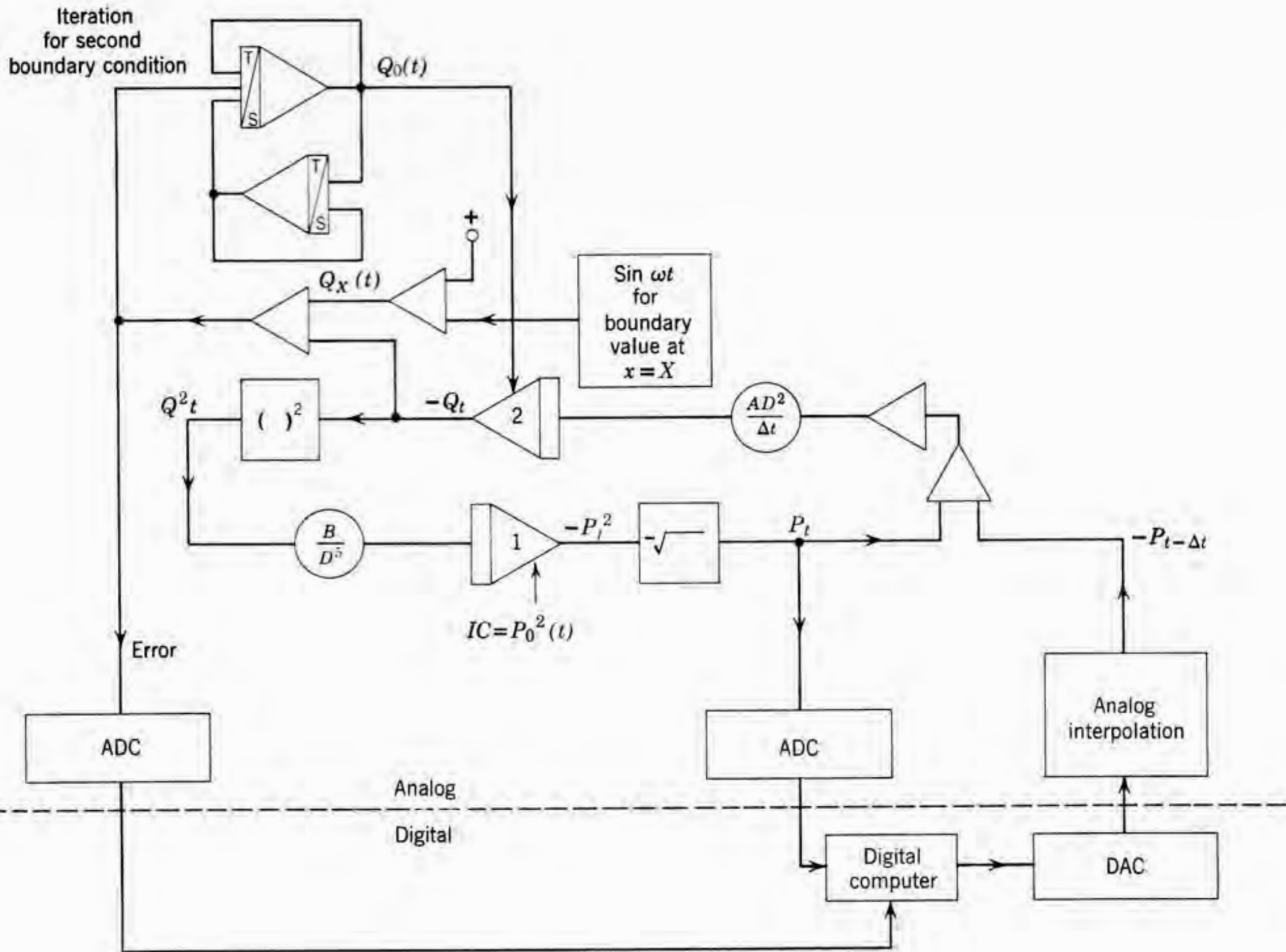
Vannevar Bush's Differential Analyser





General-Purpose Analog Computer





Hybrid Computers



Analog Time

- State evolves as time progresses
- Time is dense or continuous



Turing [1948]

- The states of 'continuous' machinery ... form a continuous manifold, and the behaviour of the machine is described by a curve on this manifold.



Baron Kelvin's Tide Predictor



Britt Phillips' Water Computer



Toy Problem: Mortar

- t time signal
- g, a, s inputs

$$x := t \cdot s \cdot \cos a$$

$$y := t \cdot s \cdot \sin a$$

$$- \frac{1}{2} \cdot g \cdot t^2$$

Distance

Height

Flow & Jump



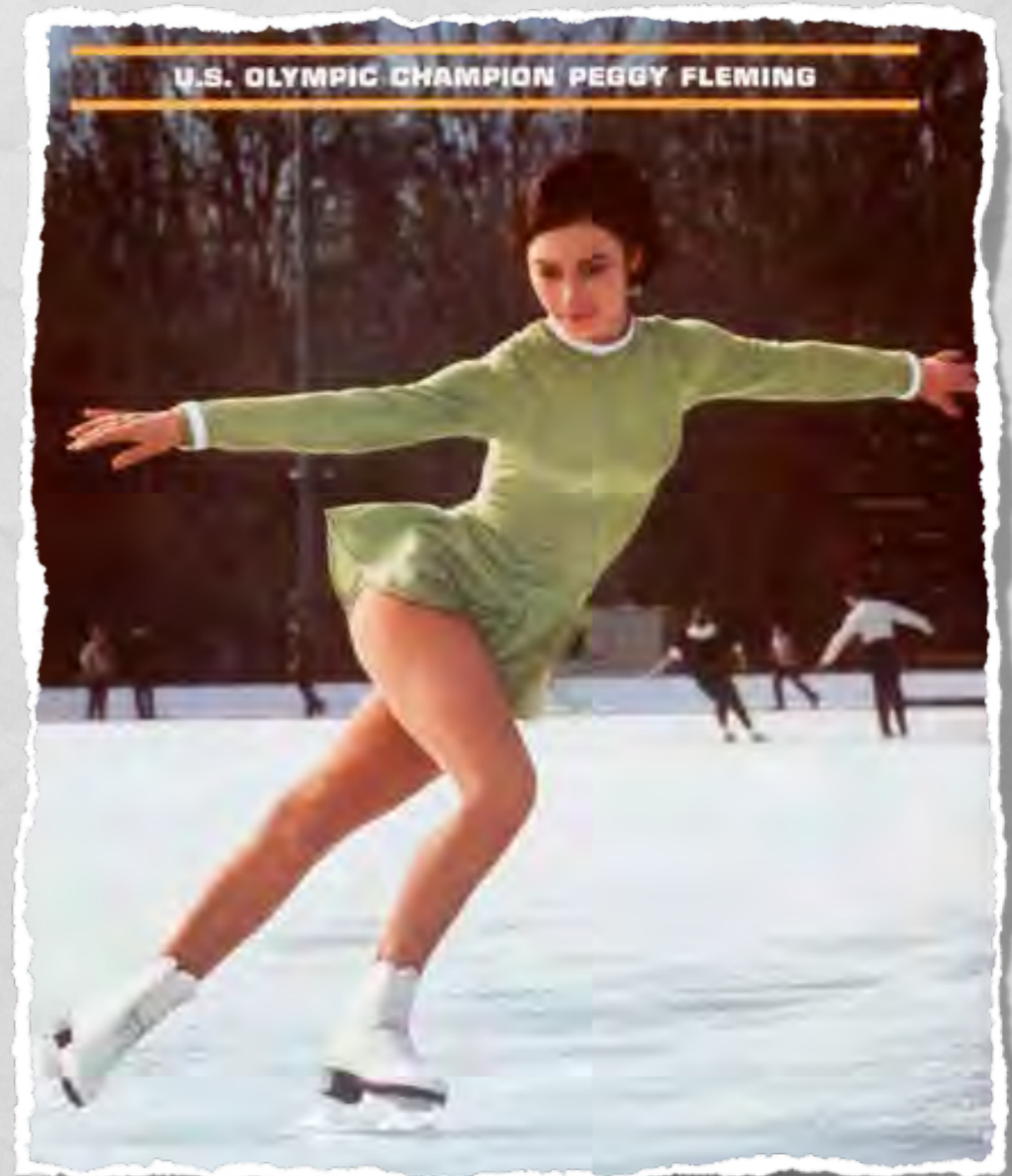
Flows

- Fixed dynamics over stretch of time



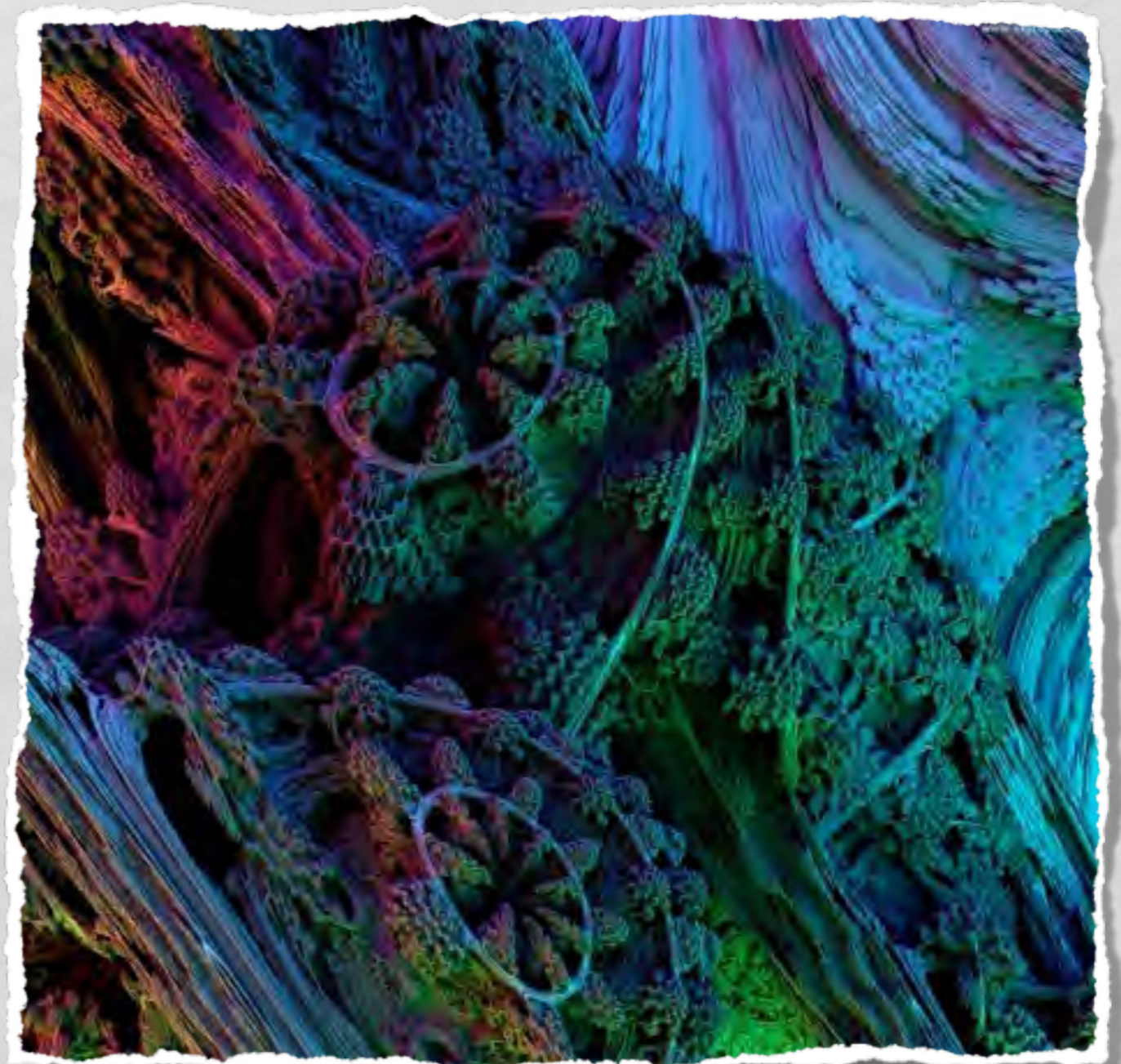
Jumps

- Change of dynamics
- Shouldn't happen too often



Sparse Jumps

- Dynamics change only finitely often in any finite trajectory



Flows

- Fixed dynamics over stretch of time
- If input wouldn't change, nothing would
- Critical equalities unchanging



A discrete algorithm is
a discrete process
whose evolution has a
finite description

An analog algorithm is
a continuous process
whose evolution has a
finite description

Algorithm

- I. An algorithm is a state-transition system
- II. Logical structures capture salient aspects of states
- III. The transition relation can be described finitely

Simple

- No input signal other than time
- Explicit (solved) equations
- Ignore output or interaction

Transition System

State

Transition

Transition System

Algorithmic

State

P
r
o
g
r
a
m

Transition

Transition System

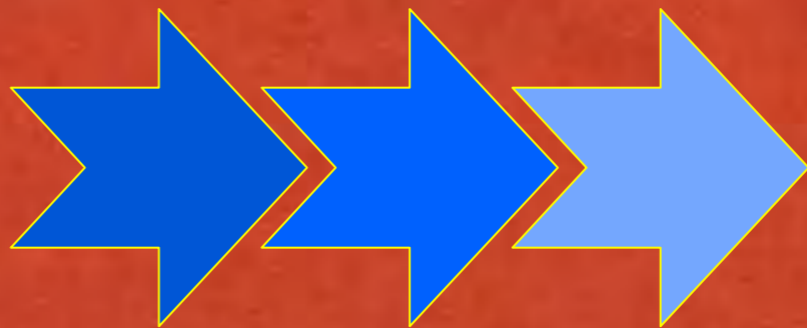
Effective

State

Initial
State

P
r
o
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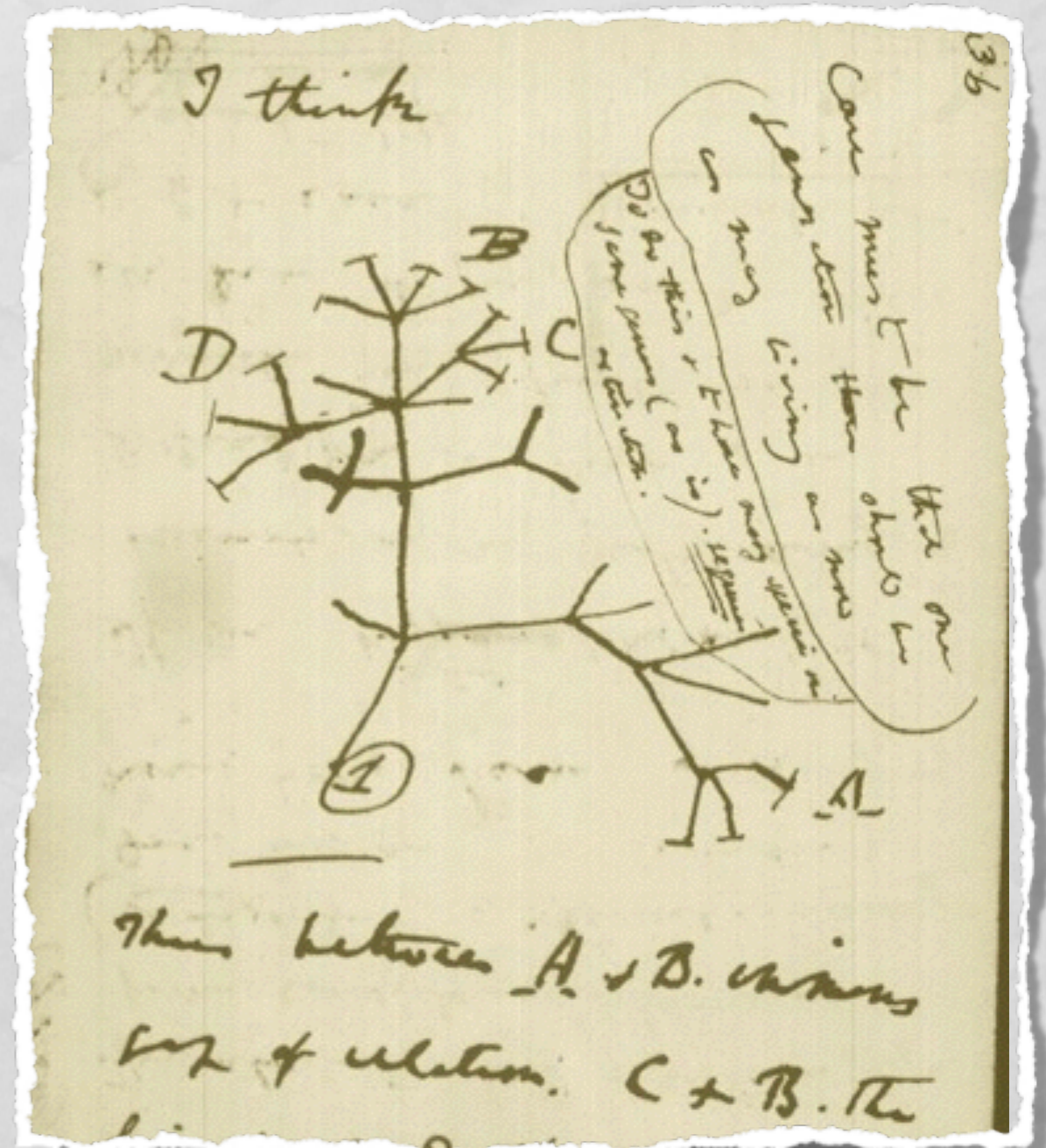
Transition



I

Timelines

- Sequential/
Branching
- Discrete/
Continuous
- Dense/Sparse
- Finite/Infinite/
Transfinite



Hartley Rogers, Jr.

For any given input, the computation is carried out in a discrete stepwise fashion, without use of continuous methods or analogue devices.

Discrete Algorithm

- A discrete state-transition system



Analog Algorithm

- A continuous state-transition system.



Discrete Transitions

My 10 Day Forecast

Updated: May 19, 2012, 11:15pm Local Time

Today
May 20



29°C 16°C

CHANCE OF RAIN: 60%
WIND: SE at 21 km/h

PM T-Storms

[Details](#)

Mon
May 21



30°C 16°C

CHANCE OF RAIN: 0%
WIND: SSE at 26 km/h

Partly Cloudy

[Details](#)

Tue
May 22



28°C 16°C

CHANCE OF RAIN: 0%
WIND: ENE at 8 km/h

Cloudy

[Details](#)

Wed
May 23



28°C 16°C

CHANCE OF RAIN: 0%
WIND: SE at 11 km/h

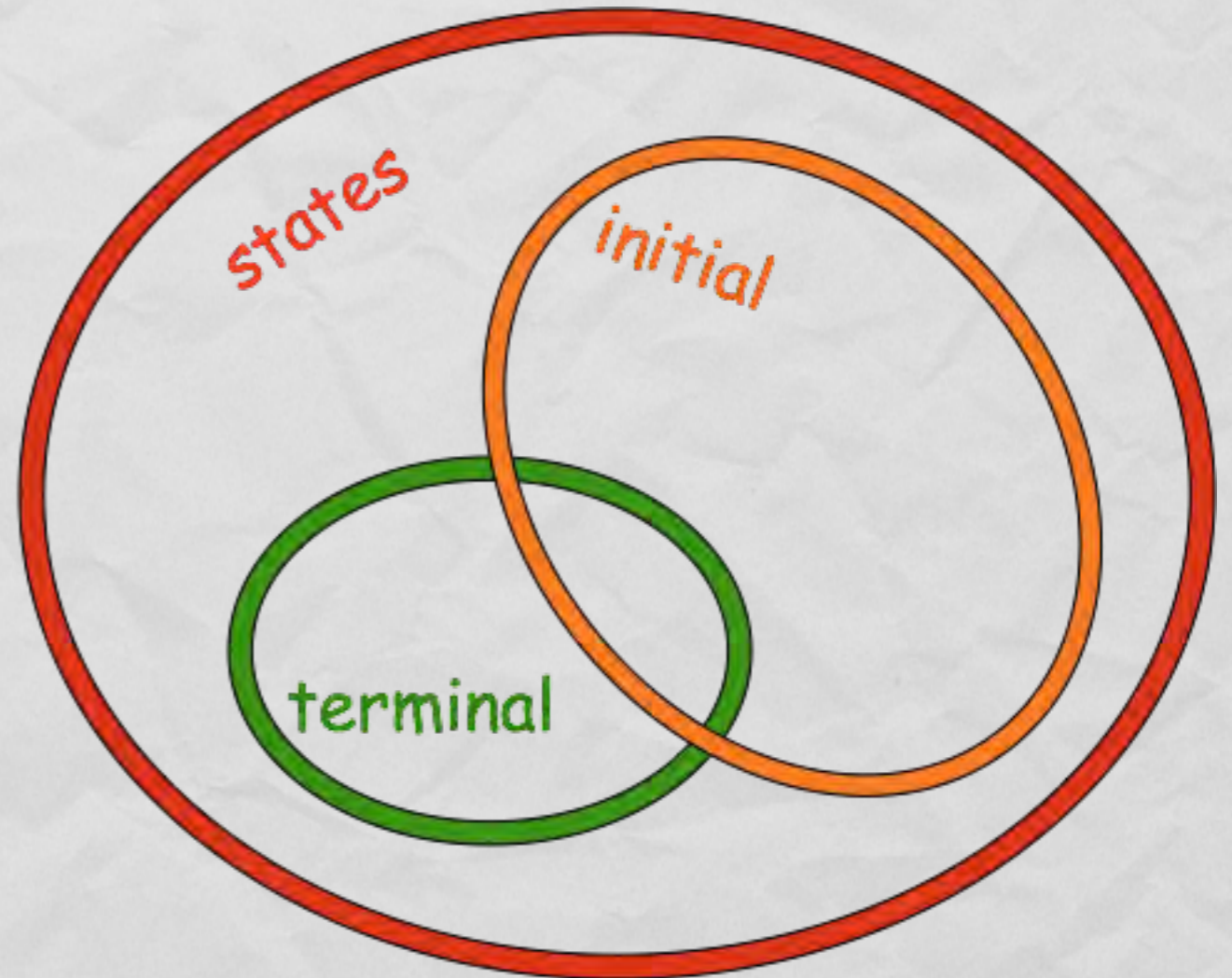
Mostly Sunny

[Details](#)

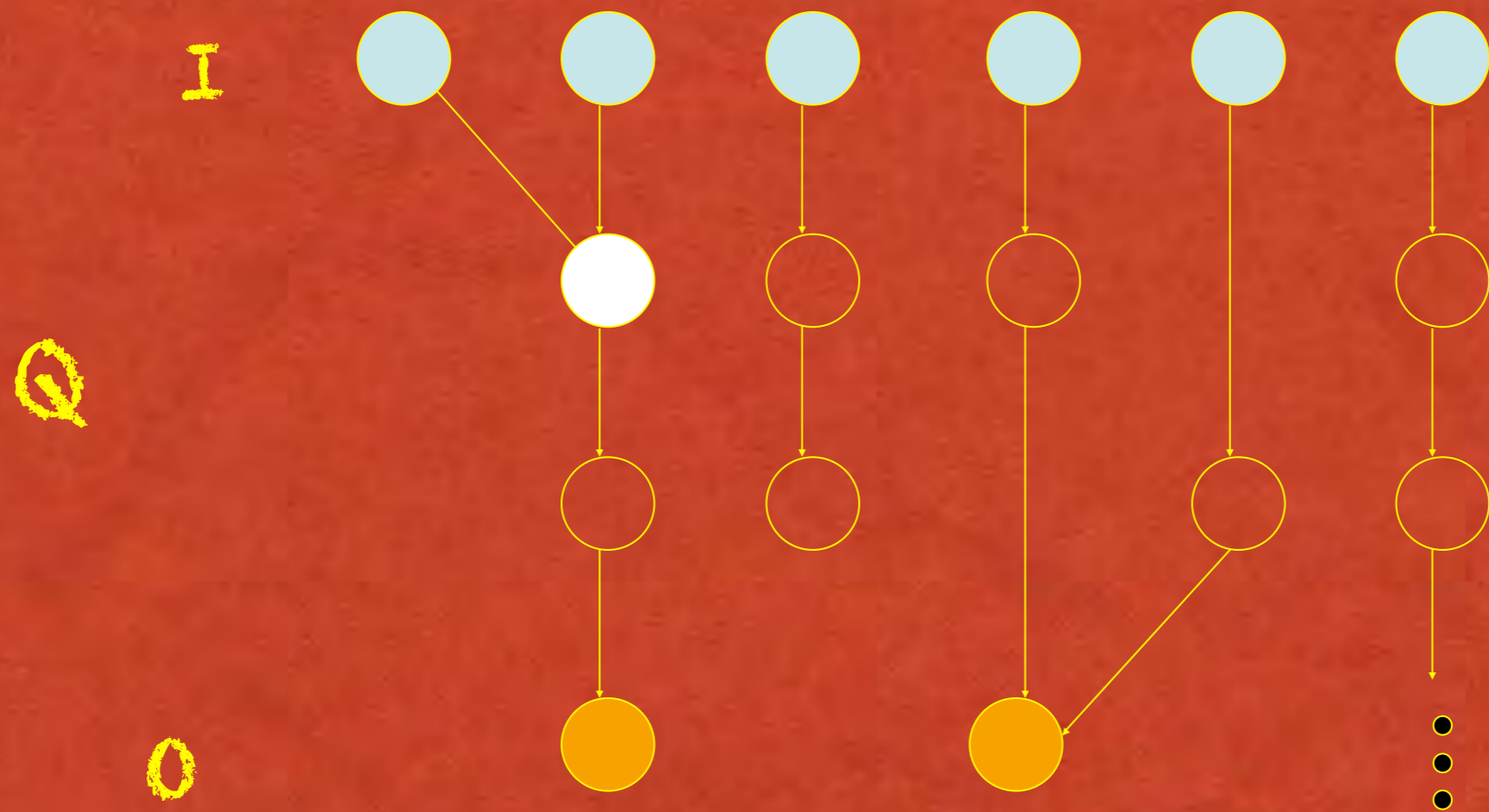


States

- Everything needed
- Initial & terminal states



Discrete Transition System



Continuous System



Infinite Change

- Zap f
- $f(t) := 0$
- Uncountably many changes



I. Evolution

- An algorithm is a state-transition system.
- Its transitions are partial functions.

Intermediate States

Ce genre de
Peinture ... de
pouvoir être
interrompu quand
on veut & repris
de même

Paul Romain Chaperon,
Traité de la peinture au
pastel (1788)



Intermediate States



Inputs

- Environment provides inputs



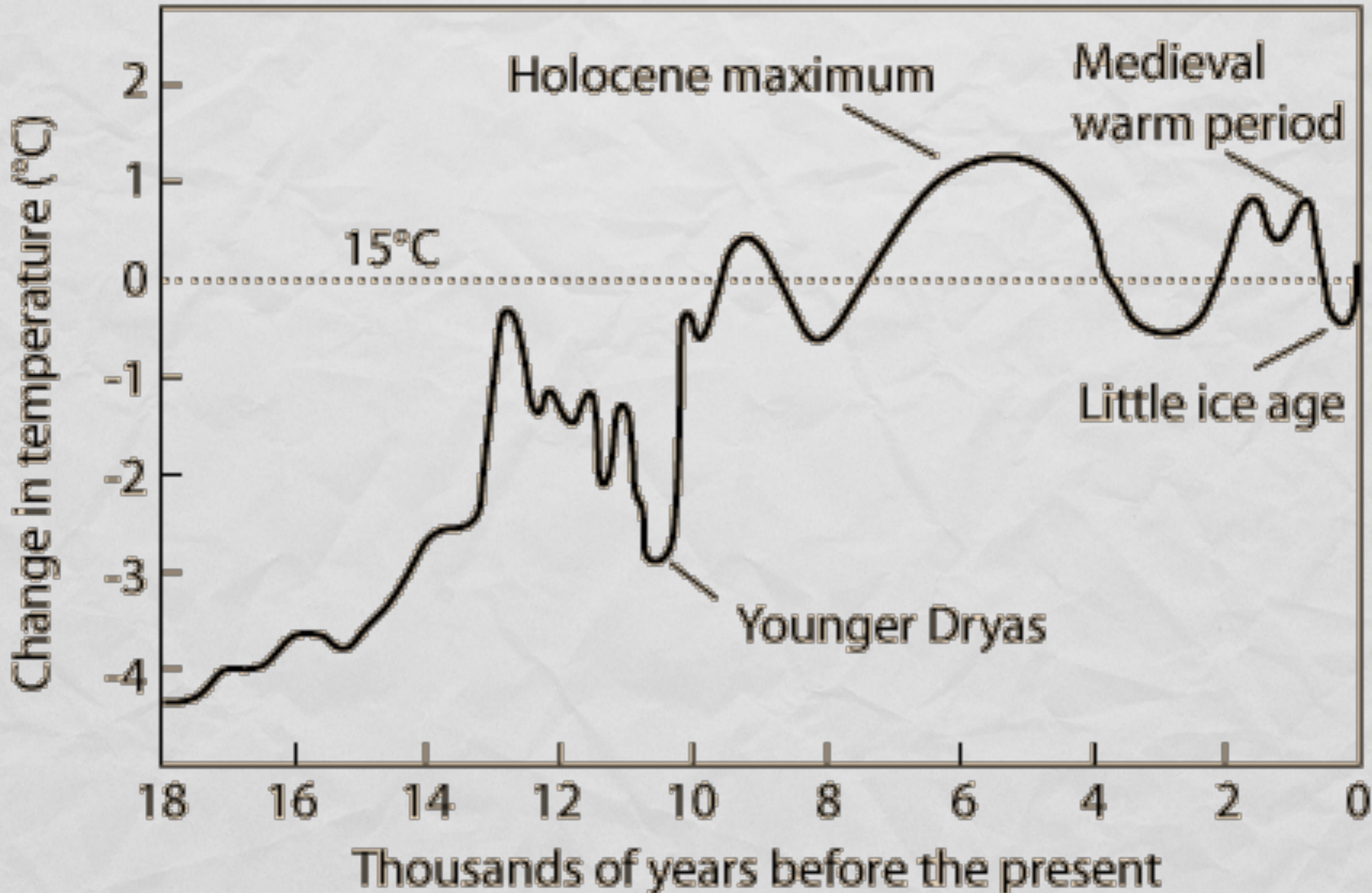
Discrete Input

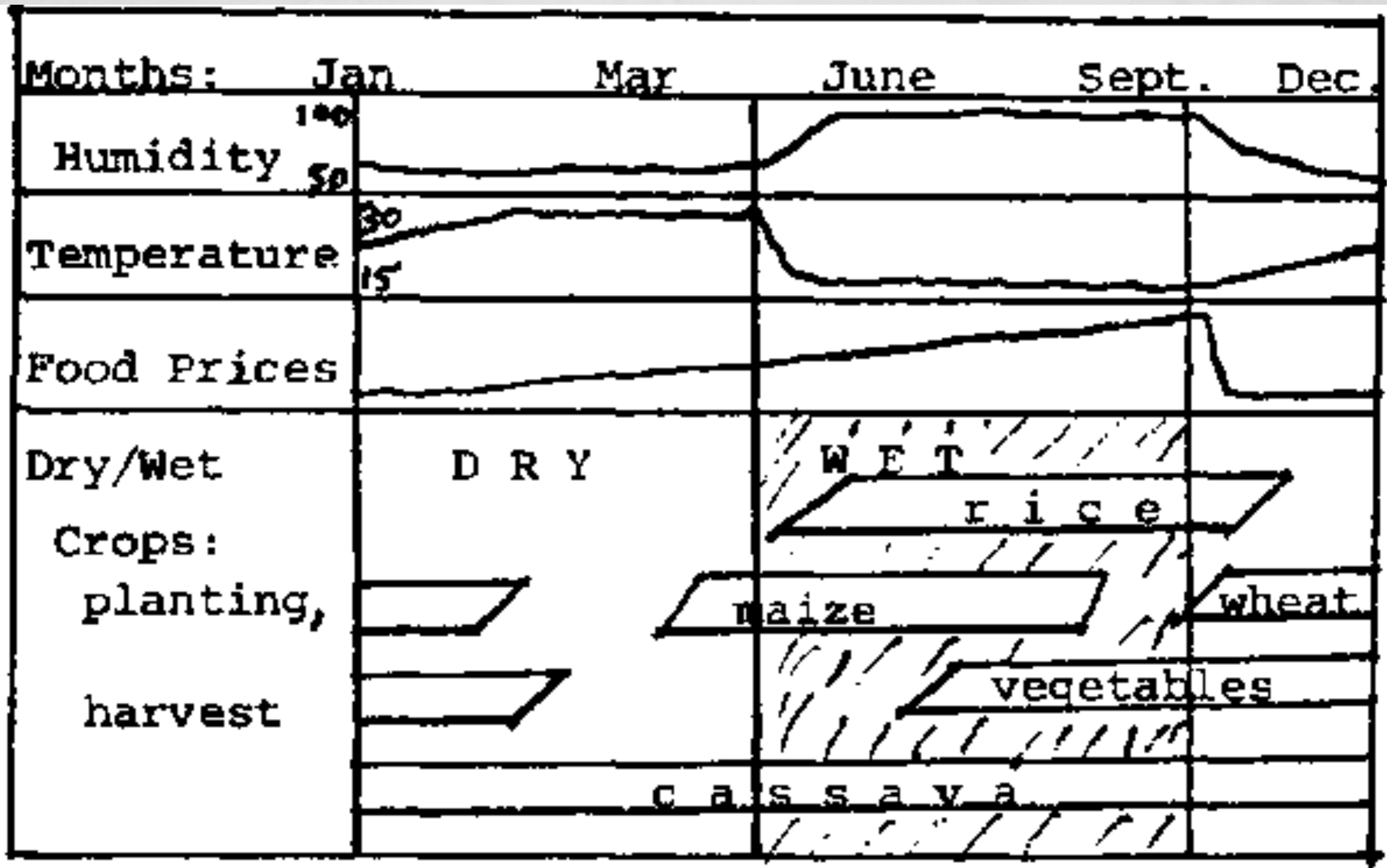
- Alternate environment steps and algorithm steps

Signals

- Function from time to domain (closed under isomorphism)
- Concatenation is associative









II



State encapsulates all relevant data!

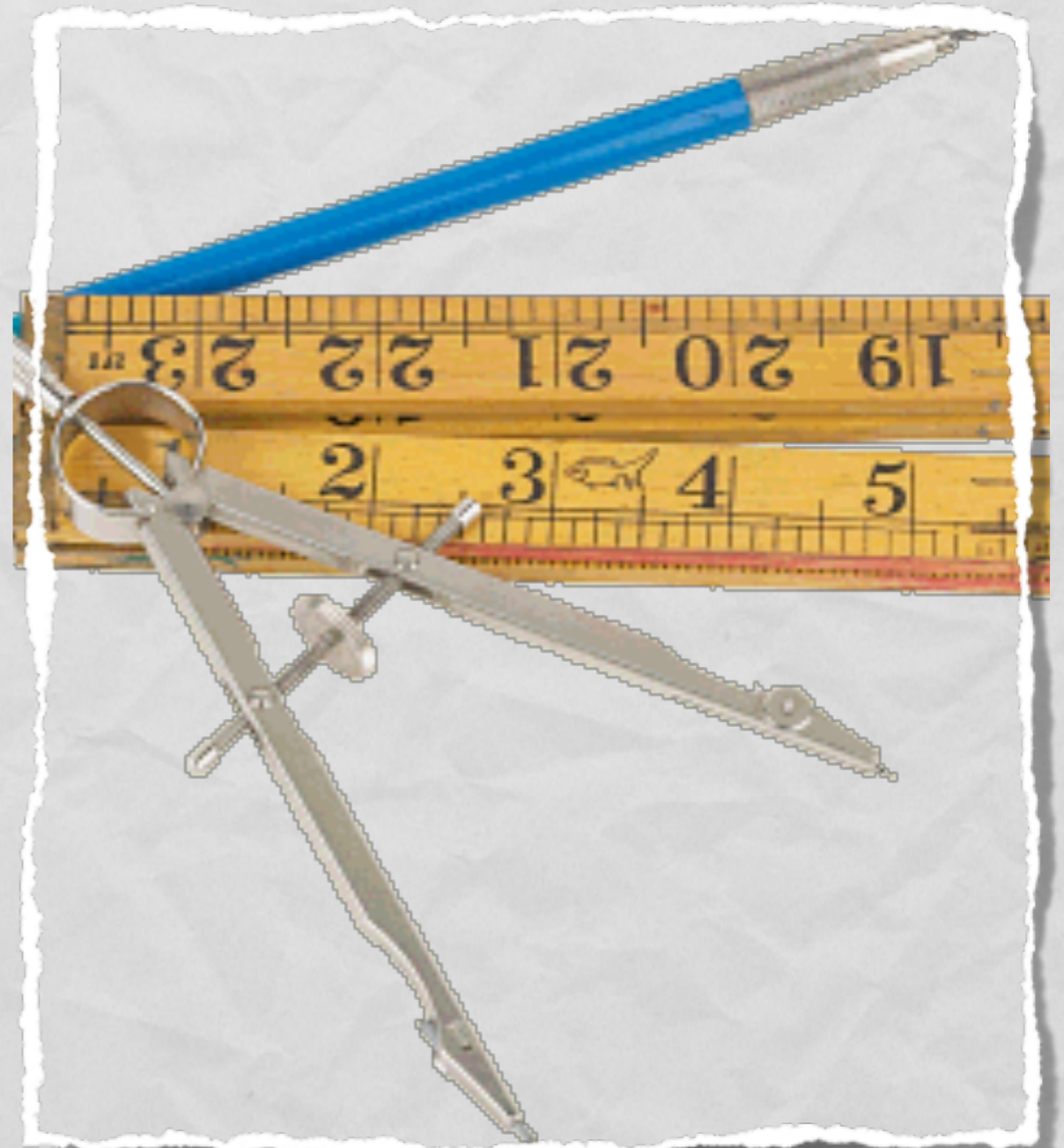
States

- Everything needed to proceed (besides the algorithm)



Geometry

- Domain
(underlying set):
points, lines, rays,
circles, tuples
and small bags
- Vocabulary &
Operations:
Compass; Ruler;
=; n ; Tuple & Bag



a b

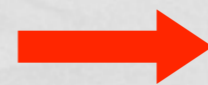
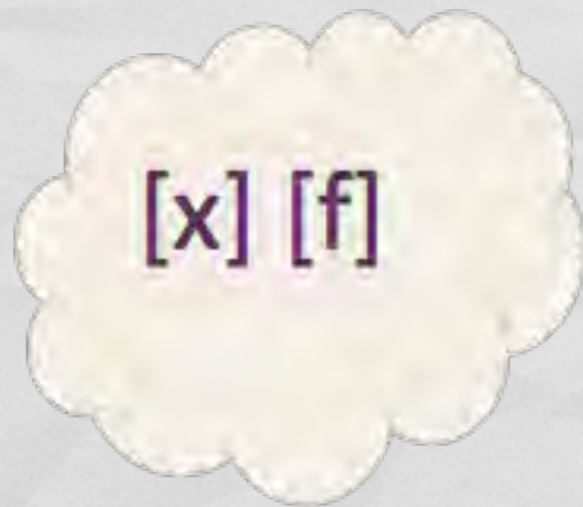
f e

+ - / sgn

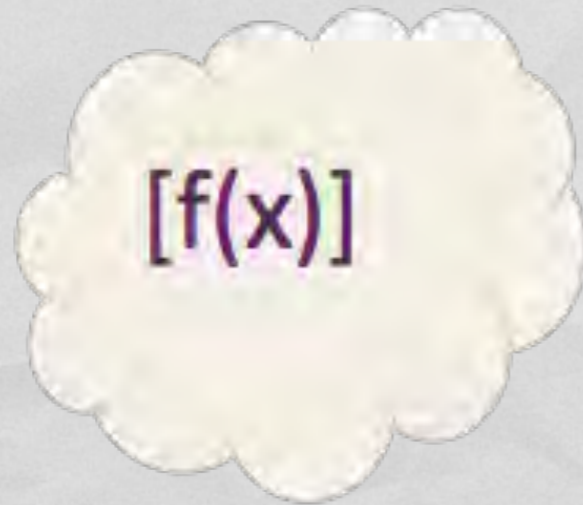
= \mathbb{F}



Algebras

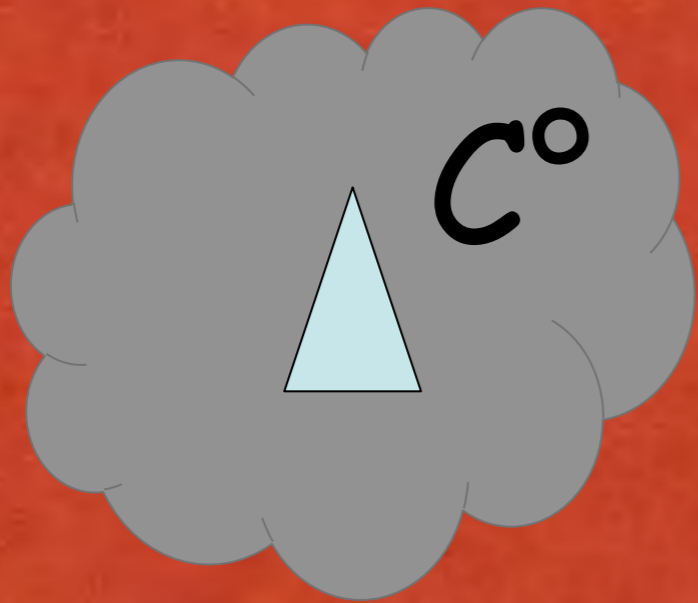


Transitions change interpretations



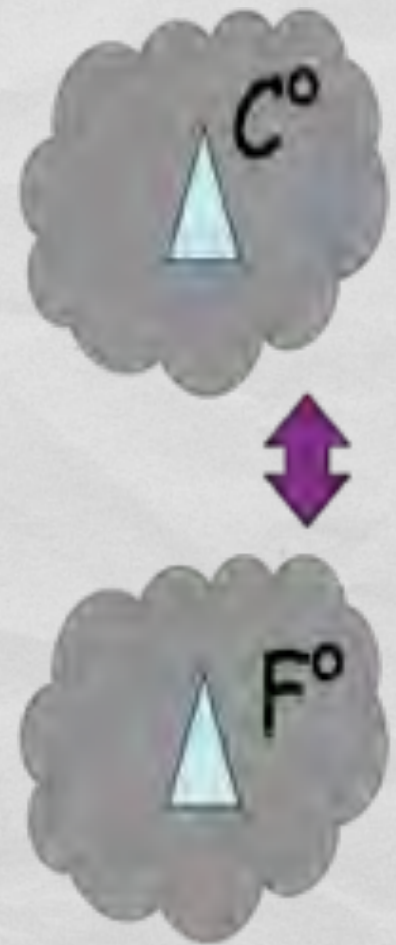
II. Abstract State

- States are (first-order) structures.
- All states share the same (finite) vocabulary.
- Transitions preserve the domain (base set) of states.

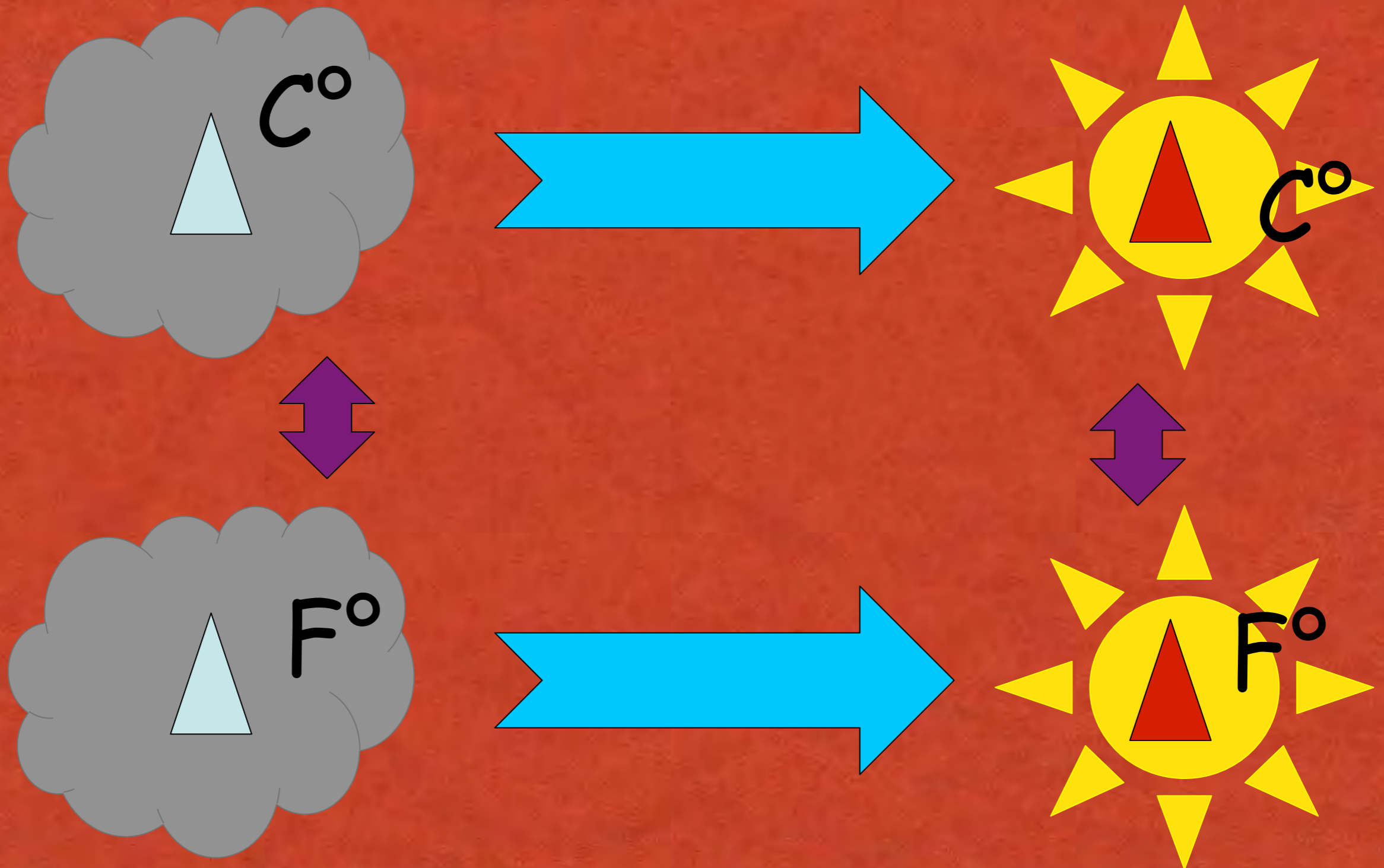


States & Transitions

- States are abstract (closed under isomorphism)
- Behavior does not depend on internal representation



Transitions respect isomorphisms



Isomorphism

- Transitions respect isomorphisms
 - $X \cong Y \Rightarrow X' \cong Y'$



Isomorphism

- Transitions respect isomorphisms
 - $X \cong Y \Rightarrow X_t \cong Y_t$



Isomorphism

- Transitions respect isomorphisms
 - $X \cong Y \ \& \ u \cong v \Rightarrow X_u \cong Y_v$



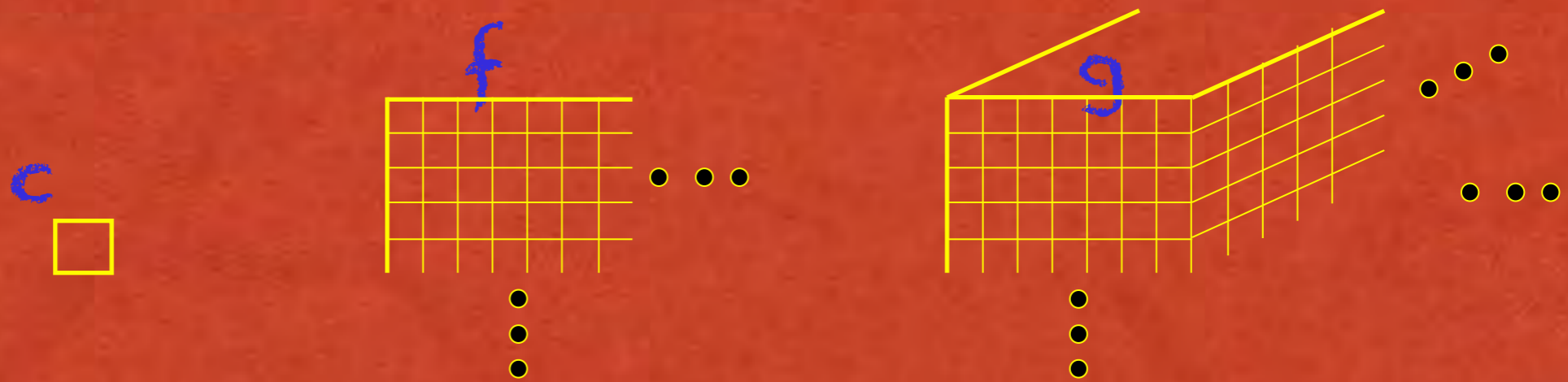
Inputs

- Environment provides inputs via ports
- X_u state after input u



States are Abstract

- Data is arranged in a structure of a finite signature.



- An algorithm is abstract, thus applicable to all isomorphic structures.



Compute by Analogy



What is Analog Computation?

- Identifying the motivating problem.
- Specifying the problem by a system of differential equations.
- Designing a network to solve the equations.
- Calculating conditions on the data and parameters to ensure good experimental behaviour of the network.
- Constructing an analog machine using a particular technology.
- Using the machine for measurements/experimental procedures.

Models

$$\frac{\partial c_1}{\partial \theta} = \frac{1}{P_v} \frac{\partial^2 c_1}{\partial Z^2} - \frac{\partial c_1}{\partial Z} - \frac{N}{M} (c_1 - c_2) \quad (1)$$

$$\frac{\partial c_2}{\partial \theta} = \frac{1}{P_{vt}} \frac{\partial^2 c_2}{\partial Z^2} + N \varepsilon (c_1 - c_2) - M f (c_2 - c_3) \quad (2)$$

$$\frac{\partial c_3}{\partial \theta} = M f (c_2 - c_3) - Q f c_3 \quad (3)$$

$$f = \exp(-P_p Z) \quad (4)$$

Signals

- Function from time to domain (closed under isomorphism)
- Concatenation is associative



Retrospection

- Current state depends on past
- Intermediate states
- $X_{uv} = (X_u)_v$



II. Abstract State

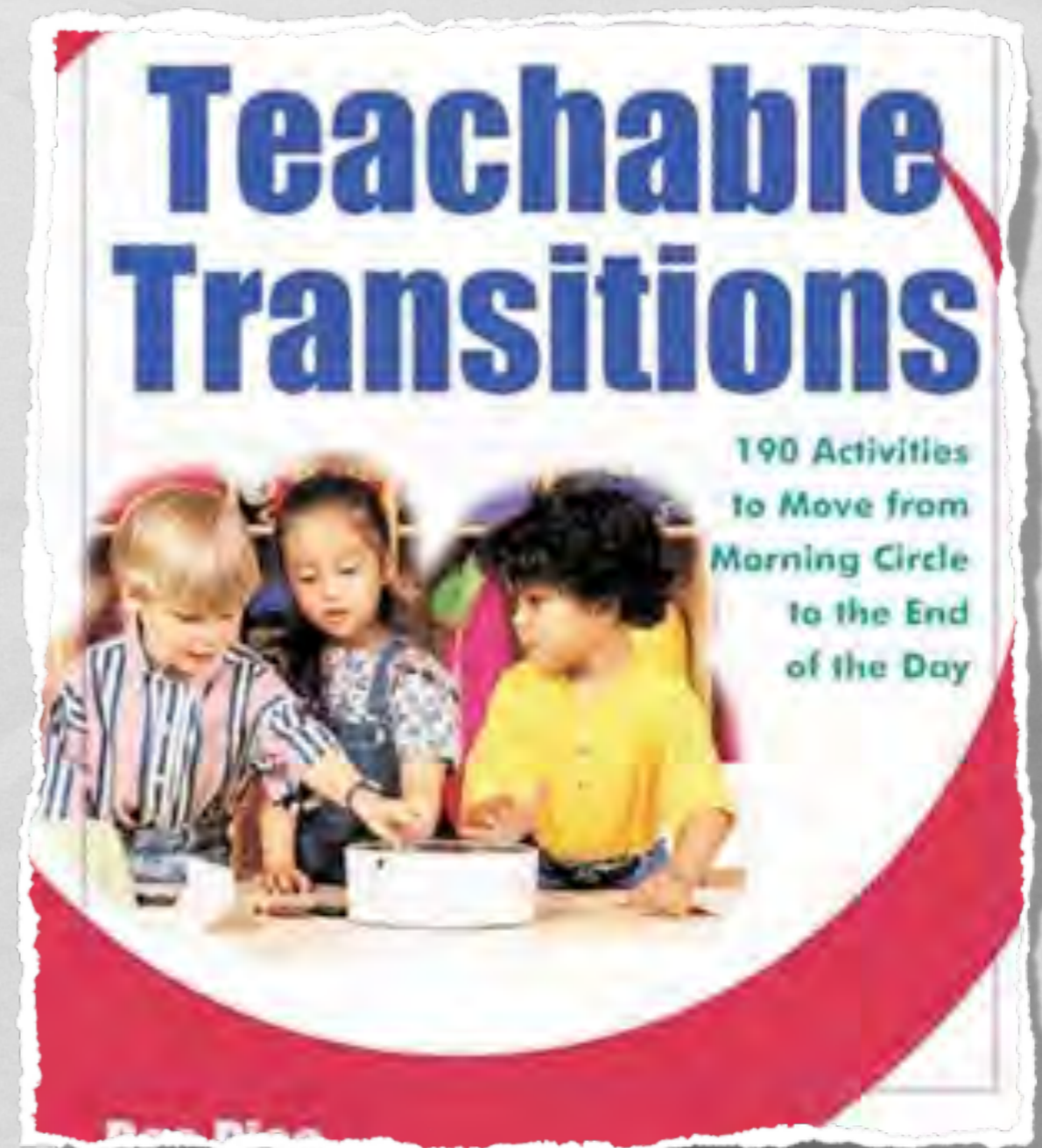
- States are (first-order) structures.
- All states share the same (finite) vocabulary.
- Transitions preserve the domain (base set) of states.
- States (and initial and terminal states) are closed under isomorphism.
- Transitions commute with isomorphisms.



III

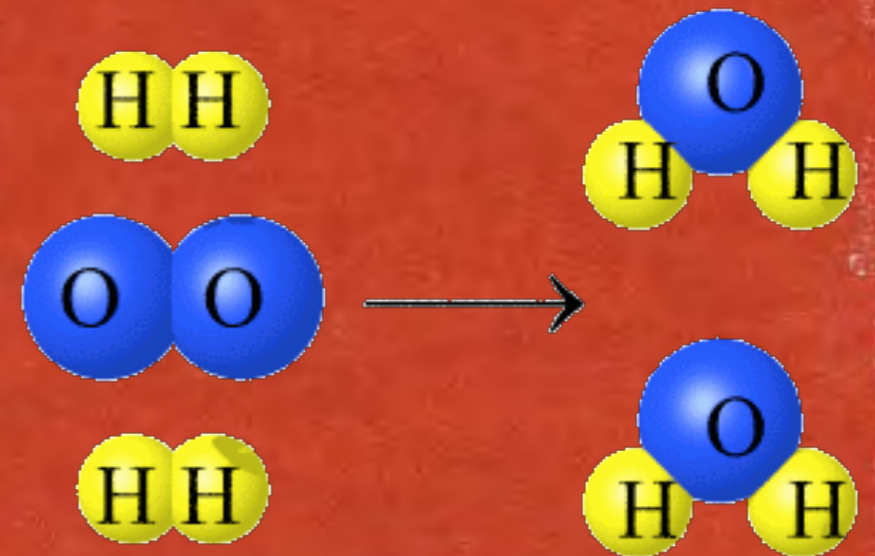
Algorithmic Transitions

- States evolve
- Evolution described by finite program



What is a transition?

- Transitions are algorithmic if they can all be described finitely (without presupposing any special knowledge).



Kleene

An algorithm in our sense must be fully and finitely described before any particular question to which it is applied is selected...

All steps must ... be predetermined and performable without any exercise of ingenuity or mathematical invention by the person doing the computing.

III. Algorithmic Transitions



- Transitions are determined by a fixed finite set of terms, such that states that agree on the values of these terms, also agree on all state changes.

Yuri Gurevich

Terms & Locations

$x, f(x)$



$x=3$

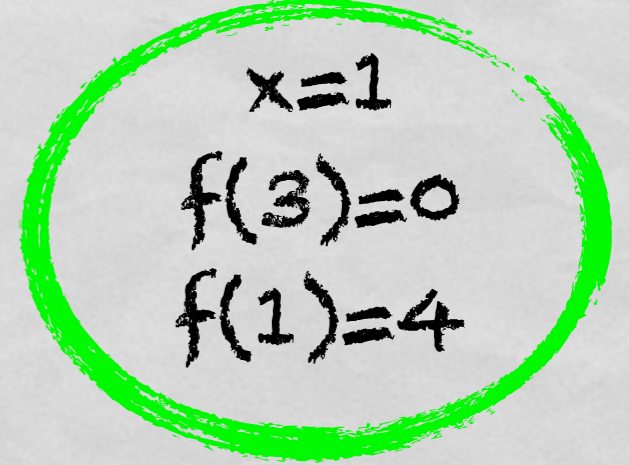
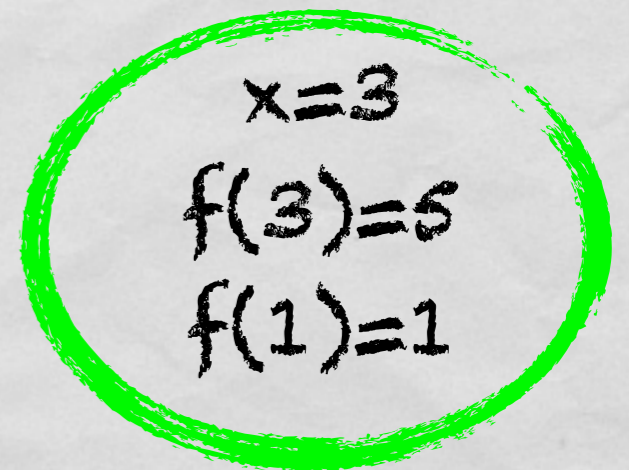
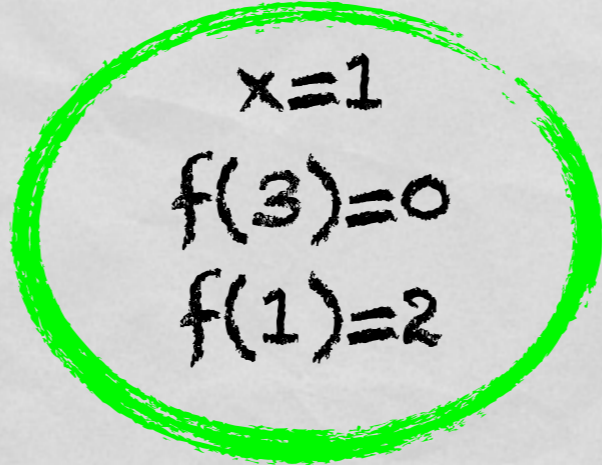
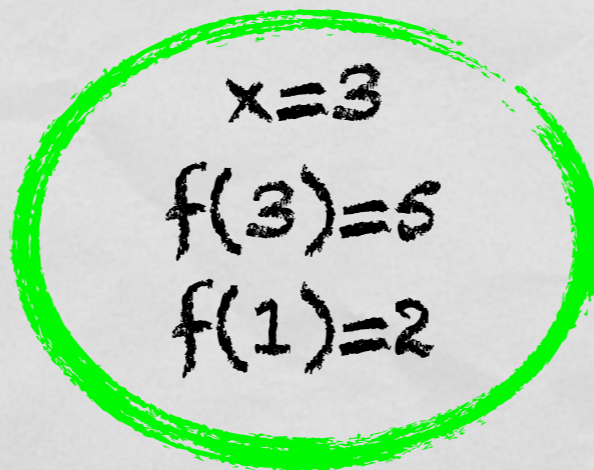
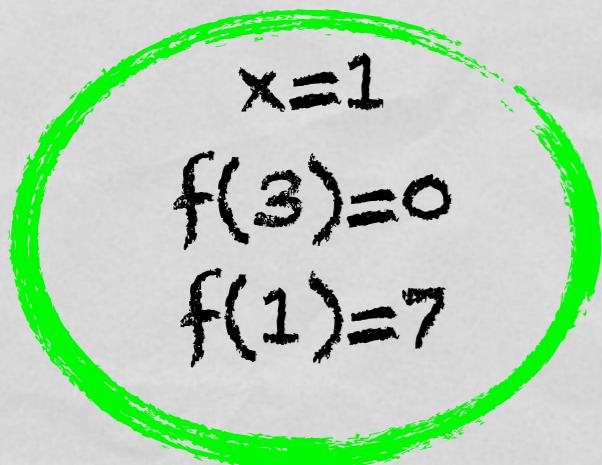
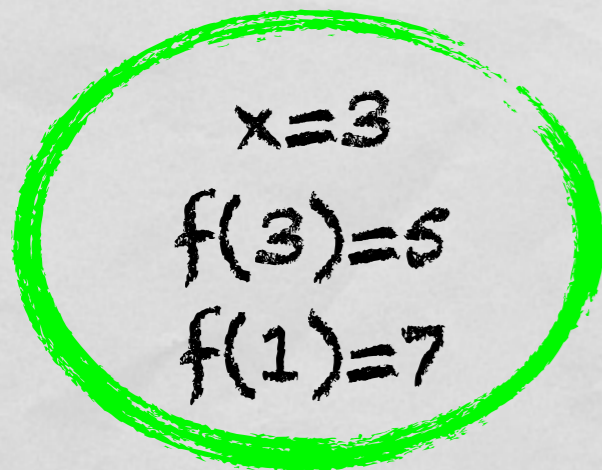


$f(3)=5$

$f(1)=2$

Critical Terms

T: $x, f(x)$

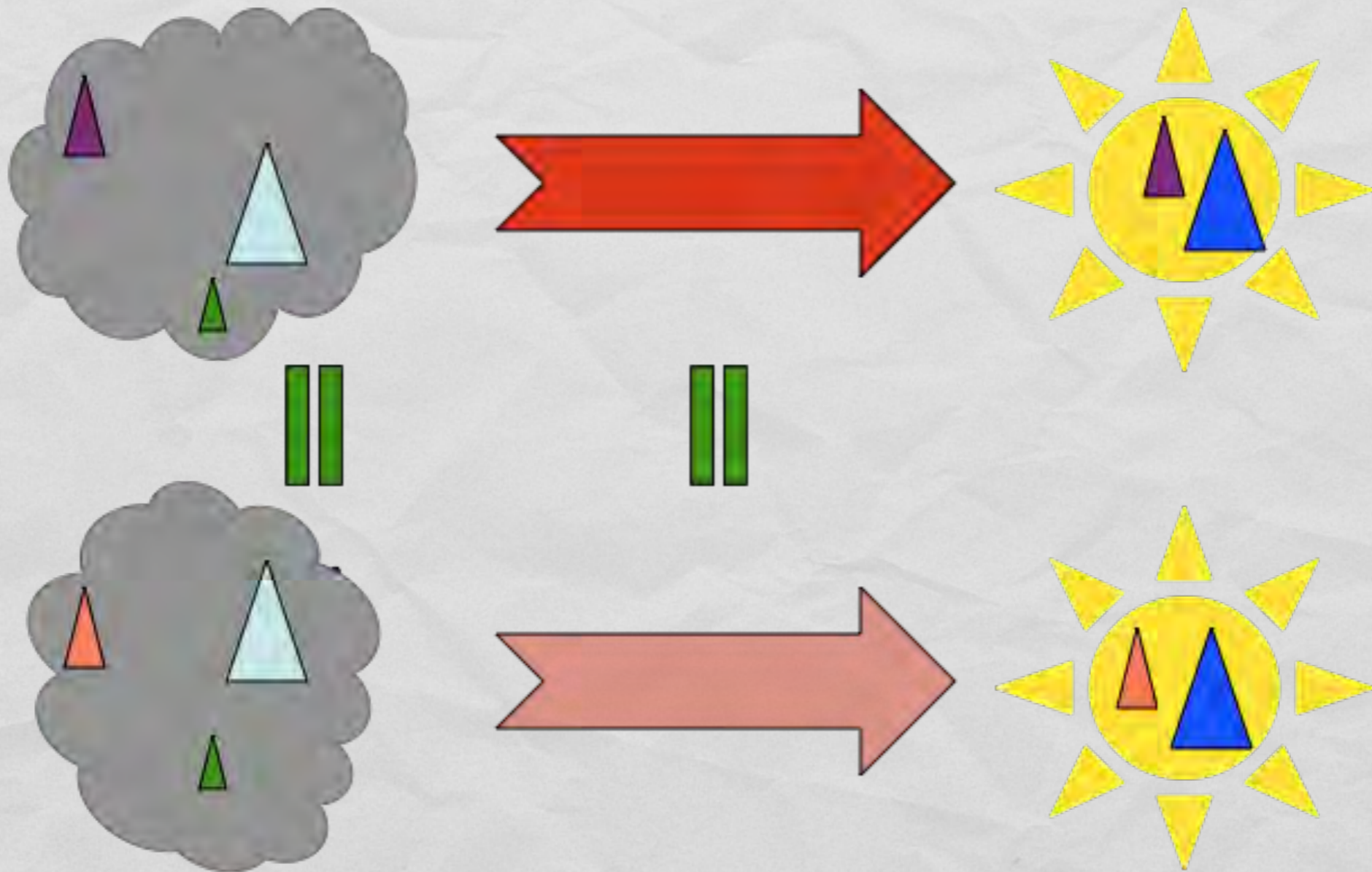


Operations

- Abstract algebraic operations
- May be partial ($3/0$ is undefined)
- Hangs when result is undefined

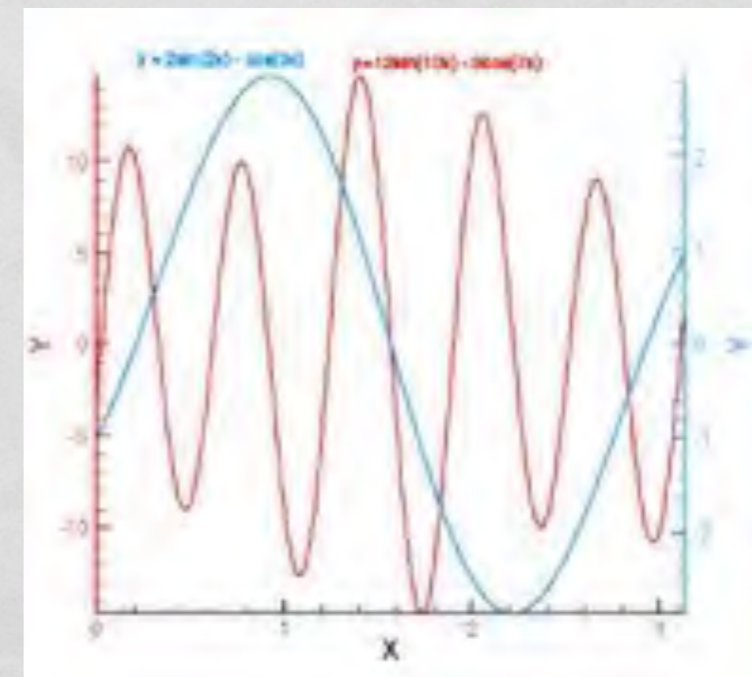


Algorithmic Transitions



Transitions

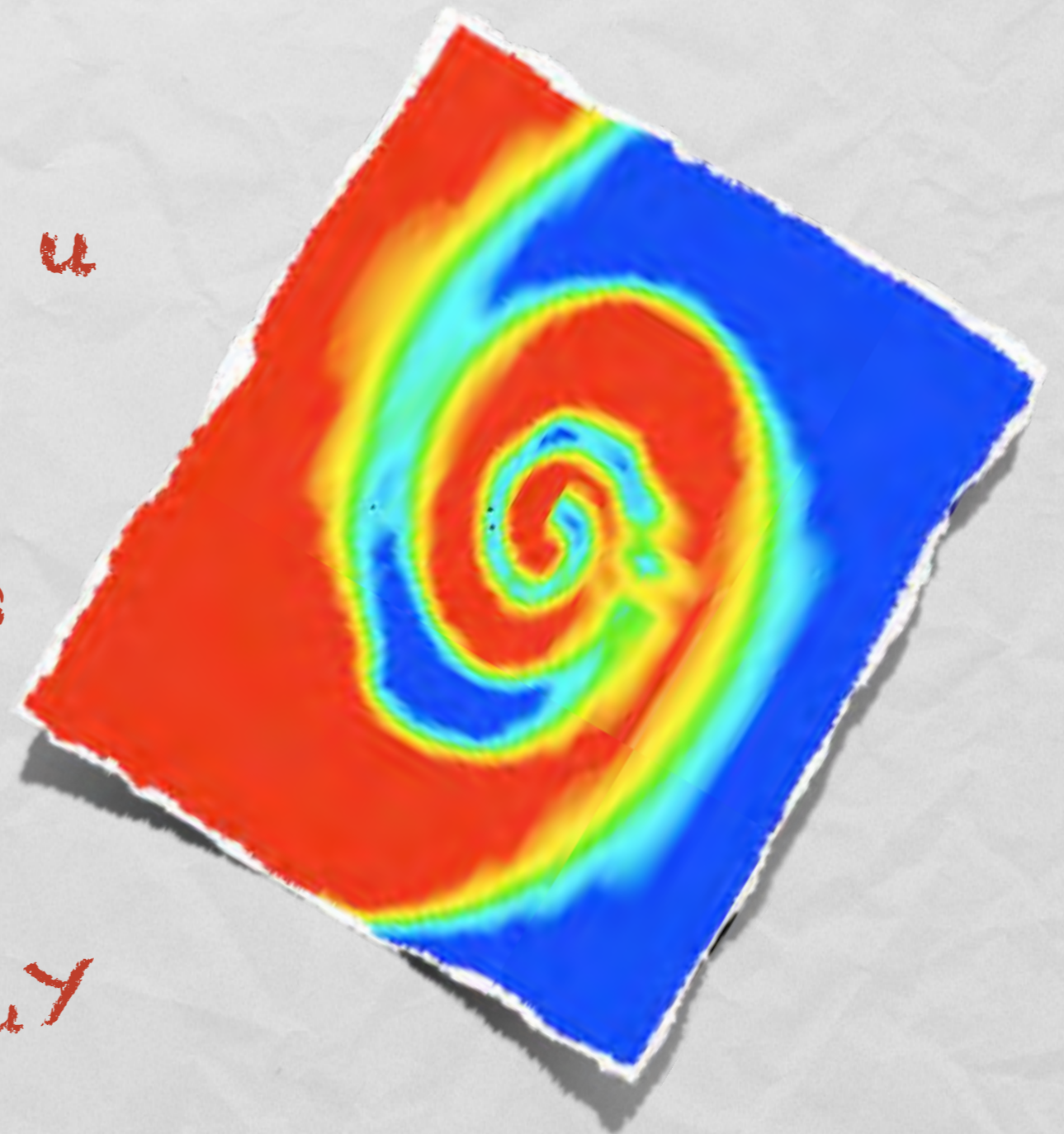
- View state as location-value pairs
- $f(a,b,c) \mapsto d$
- Changes to state X
 - $\Delta X = X' \setminus X$
 - $\Delta_u X = X_u \setminus X$



CHANGE

Evolution

- Transition to X_u under input signal u
 - $\Delta_u X = X_u \setminus X$
- Evolution depends on critical terms and input port
 - $X \stackrel{T}{=} Y \Rightarrow \Delta_u X = \Delta_u Y$



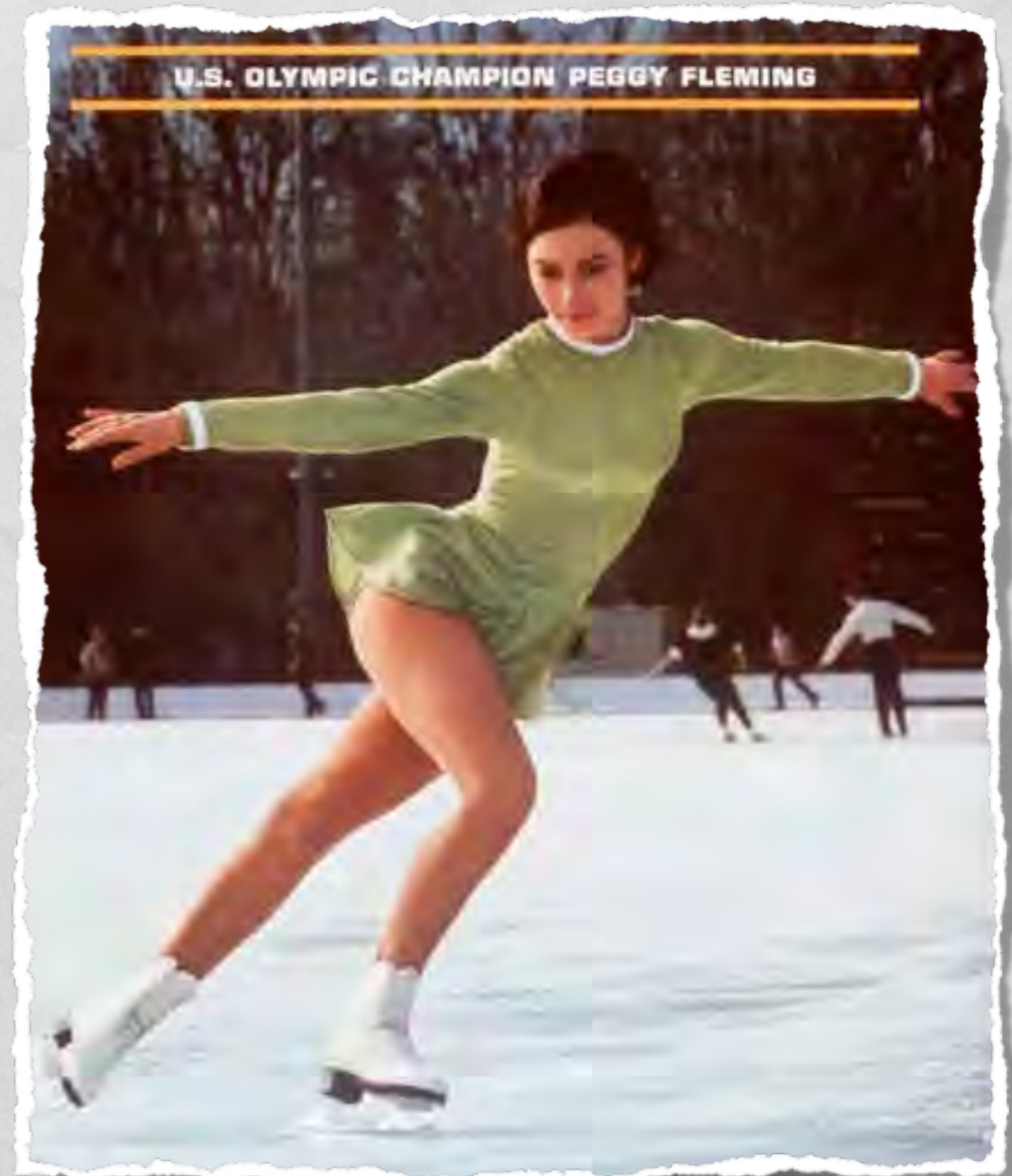
Flows

- Fixed dynamics over stretch of time
- If input wouldn't change, nothing would
- Equalities between critical terms maintained



Jumps

- Change of dynamics
- Requires conditionals



Partiality

- Locations still to be accessed are determined by locations already accessed



Past Determines Future

$$X_u =_T Y_v$$

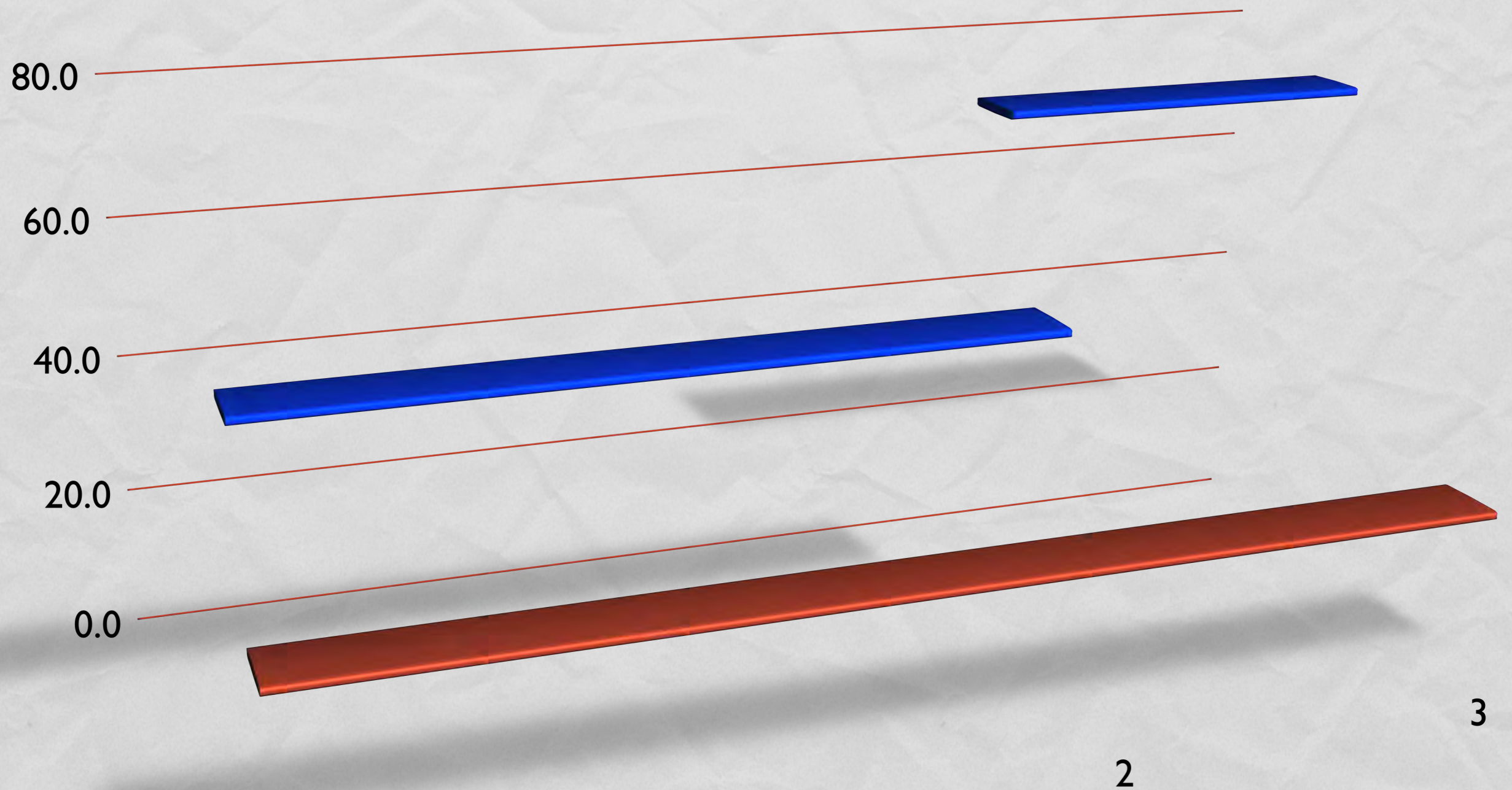
$$u|u| = v|v|$$



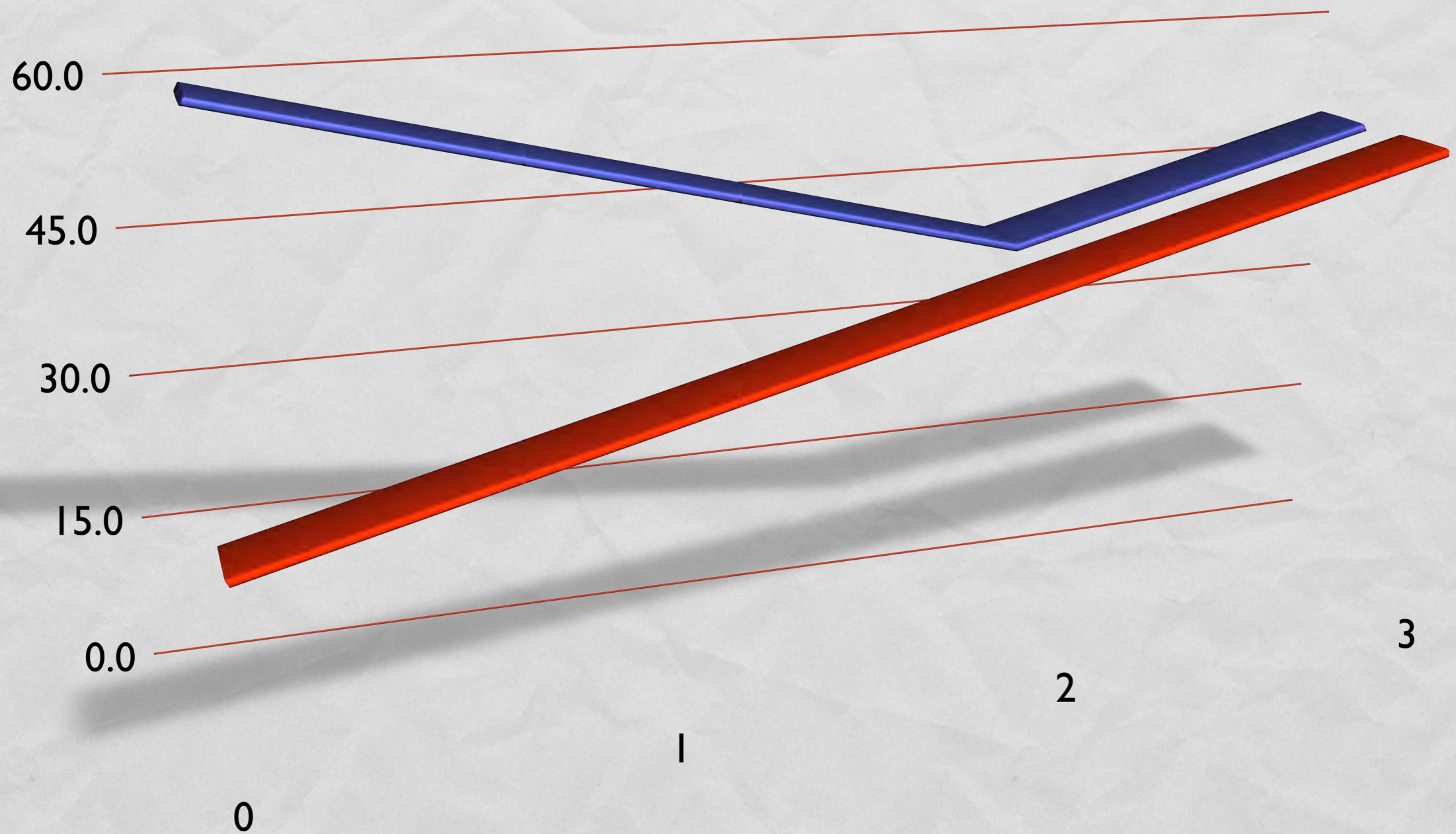
$$\Delta_u X = \Delta_v Y$$



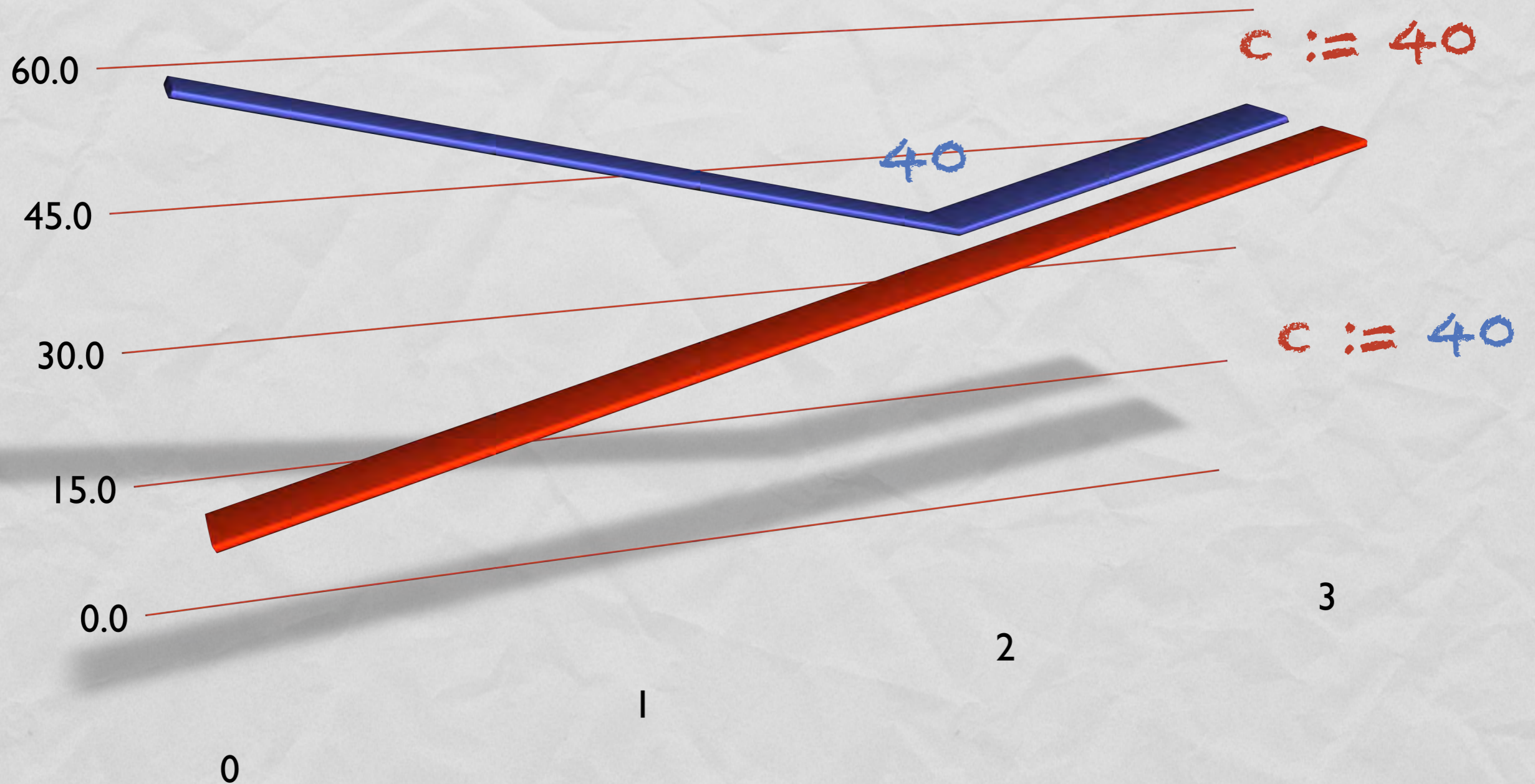
Non-Deterministic



Deterministic



Critical Moment



Non-Flows

- $x := x+1$
- if $x=0$ then $y := -y$
 - unless $y=0$
- $x := x/2$
 - unless $x=0$



Flow

vs.

Jump

$$z = \lfloor t \rfloor$$

if $t < z$ then $z := z - 1$ ||

if $t \geq z + 1$ then $z := z + 1$

Before & After



- At every jump, there are "before" and "after" states: X before algorithm makes changes and X' after

Critical Moment

- Suppose $f(a,b,c) := d$ in $\Delta_u X$
- After some prefix w of u , d in $[[T]]_{X_w}$
- If not, let Y be X with d' instead of d
- And v be u with d' instead
- By criticality $f(a,b,c) = d$ also in Y_u
- By isomorphism $f(a,b,c) = d'$ in Y_v
- So $u \neq v$, and signals must part ways

Flow Equations

- Solved form
- Left-differentiation operator
- Solver for implicit equations
- Infinitesimals

And God said:

$$\nabla \cdot \mathbf{E} = \frac{\rho}{\epsilon_0}$$

$$\nabla \cdot \mathbf{B} = 0$$

$$\nabla \times \mathbf{E} = -\frac{\partial \mathbf{B}}{\partial t}$$

$$\nabla \times \mathbf{B} = \mu_0 \mathbf{J} + \mu_0 \epsilon_0 \frac{\partial \mathbf{E}}{\partial t}$$

and there was

light! ...

Bounce, Bounce

- t time signal
- dt infinitesimal
- g, k constants
- $s := s + g \cdot dt$
- $x := x + s \cdot dt$
- if $x=0$ then $s := -k \cdot s$



Explicit Flow

- t time signal
- g, a, s inputs



$$x := t \cdot s \cdot \cos a$$

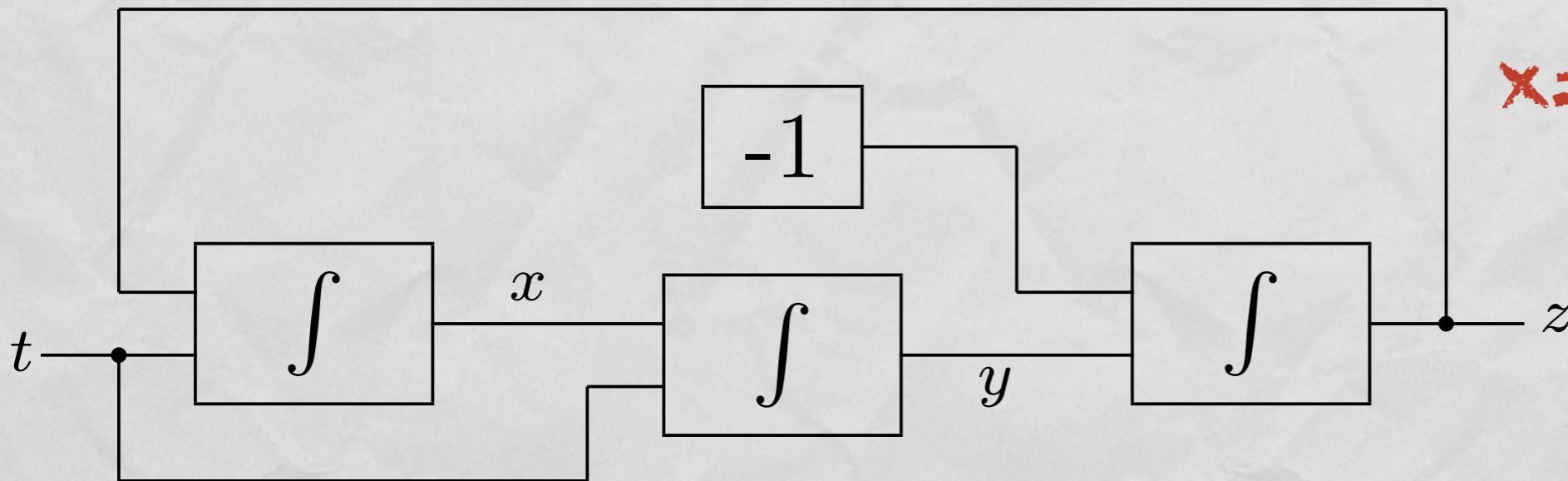
$$y := t \cdot s \cdot \sin a$$

$$- \frac{1}{2} \cdot g \cdot t^2$$

Implicit Flow

- t time signal

- initially
 $x=1; y=z=0$



$$x' = z$$

$$y' = x$$

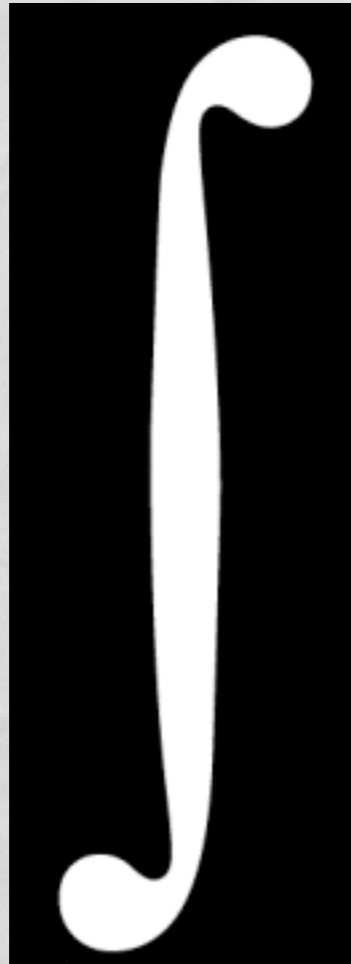
$$z' = -y$$

Signal of Signals

- Tiny pieces of history
- Left-differentiation operator
- $y := x'$



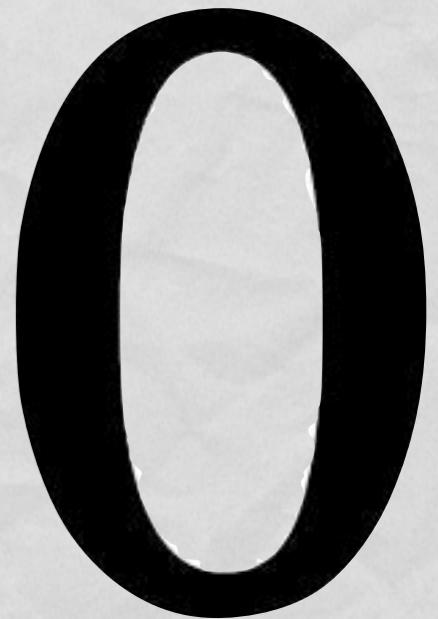
Bruno Scarpellini 1963



- Compute the undecidable via infinite precision integral.
- It is conceivable that the mathematics of a collection of axioms may lead to undecidable propositions. (2003)

John Myhill

- To make Scarpellini's work a basis for constructing an actual computer which can solve problems which are not digitally (= recursively) solvable:
- Assume perfect functioning.
- Assume perfect sensor (zero test).



Tide-Prediction Manual

- The machine to be described here, like almost every contrivance, apparatus, or machine in practical use, is based very largely upon what has been accomplished by others who previously labored in the same field.

U.S. Coast and Geodetic Survey
(1915)

IV

An algorithm is effective if its initial states have a finite description

If initial states can be
described, then all
states can be